Computer Architecture

Lecture 07 – Address Translation & Virtual Mem

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Any problem in computer science can be solved with another layer of indirection.

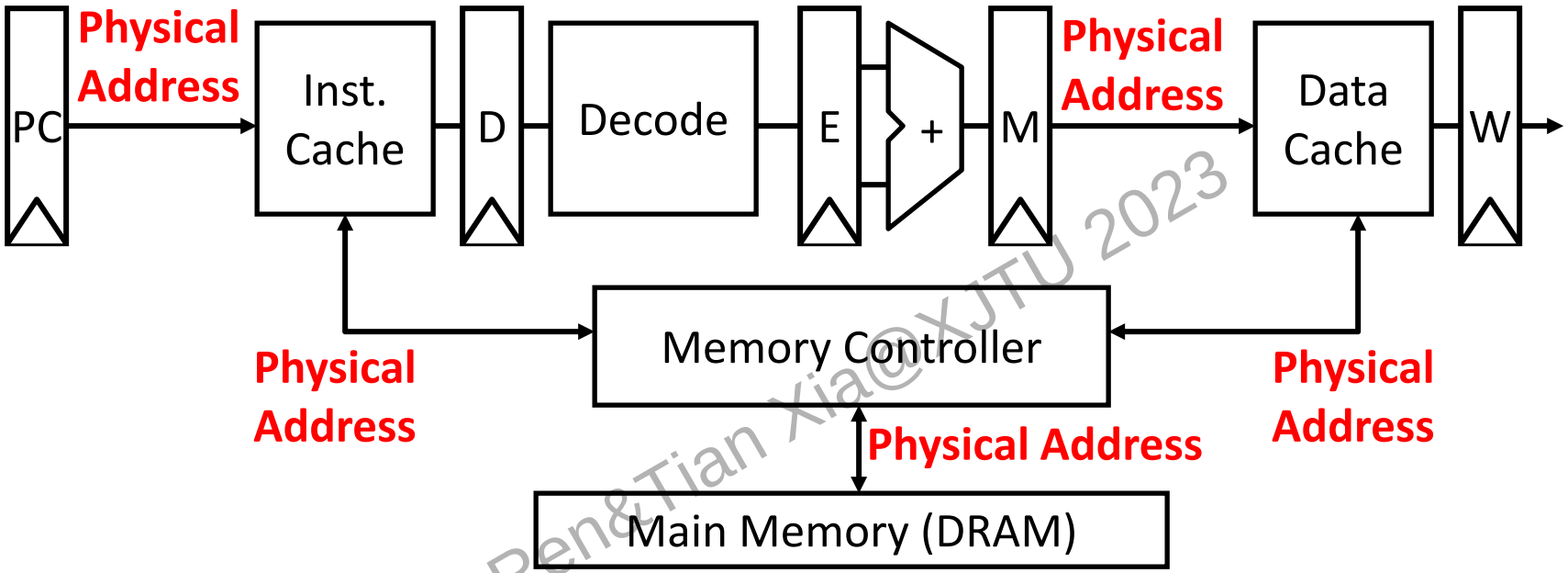
--David Wheeler

British Computer Scientist

(1927--2004)

- Fellow of the Royal Society (1981)
- Computer Pioneer Award (1985)
- Fellow, Computer History Museum (2003)

Bare Machine



In a bare machine, the only kind of address is a physical address, corresponding to address lines of actual hardware memory.

Managing Memory in Bare Machines

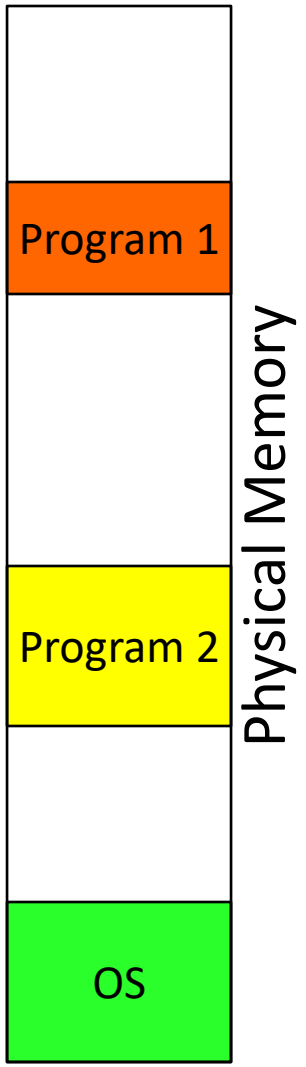
- Early machines only ran **one program at a time**, with this program having unrestricted access to all memory and all I/O devices
  - This simple memory management model was also used in turn by the first minicomputer and first microcomputer systems
- Subroutine **libraries** became popular, were written in **location-independent** form
  - Different programs use different combination of routines
- To run program on bare machines, use **linker** or **loader** program to relocate library modules to actual locations in physical memory

Dynamic Address Translation

- Motivation
  - Each process limited to a non-overlapping contiguous physical memory region (space doesn't start from addr 0…)
  - Everything must fit in the region
- Location-independent programs
  - Programming and storage management ease

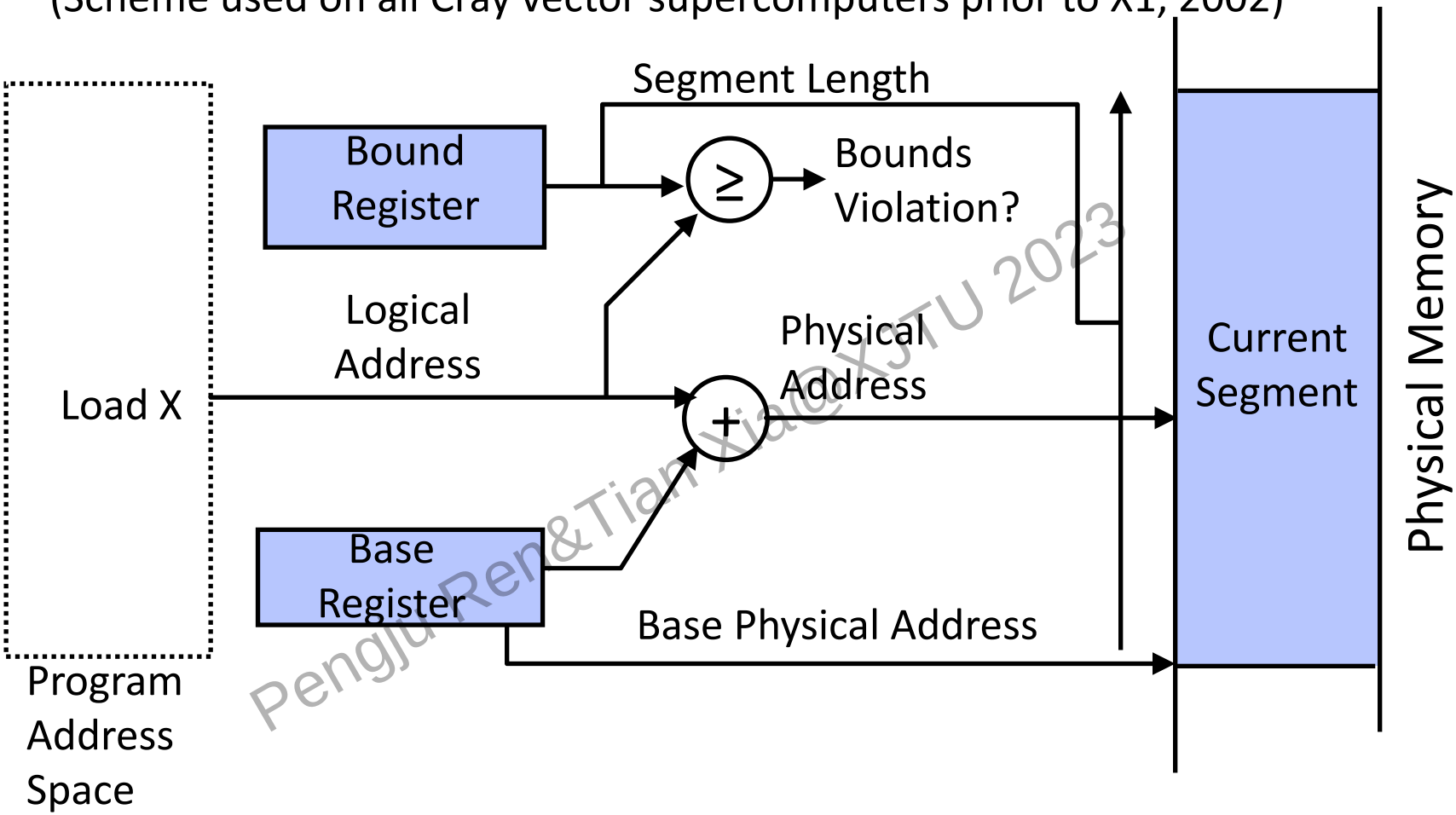
  → **need for a *base* register**
- Protection
  - Independent programs should not affect each other inadvertently

  → **need for a *bound* register**
- Multiprogramming drives requirement for **resident supervisor software** (e.g. OS) to manage context switches between multiple programs



Simple Base and Bound Translation

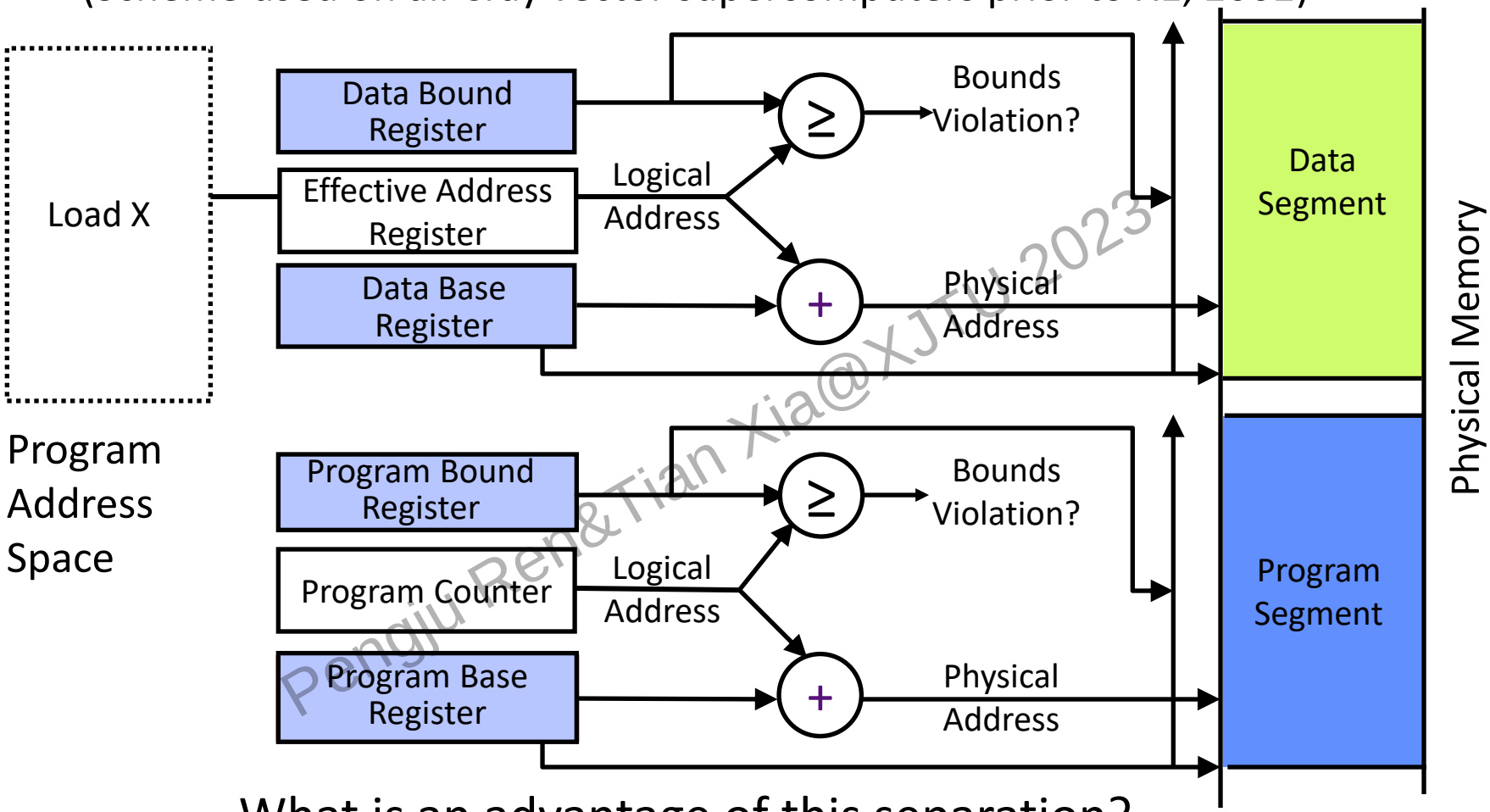
(Scheme used on all Cray vector supercomputers prior to X1, 2002)



Base(Staring address) and **bounds**(size of region) registers are visible/accessible only when processor is running in the ***supervisor mode*** *(privileged control registers)*

Separate Areas for Program and Data

(Scheme used on all Cray vector supercomputers prior to X1, 2002)

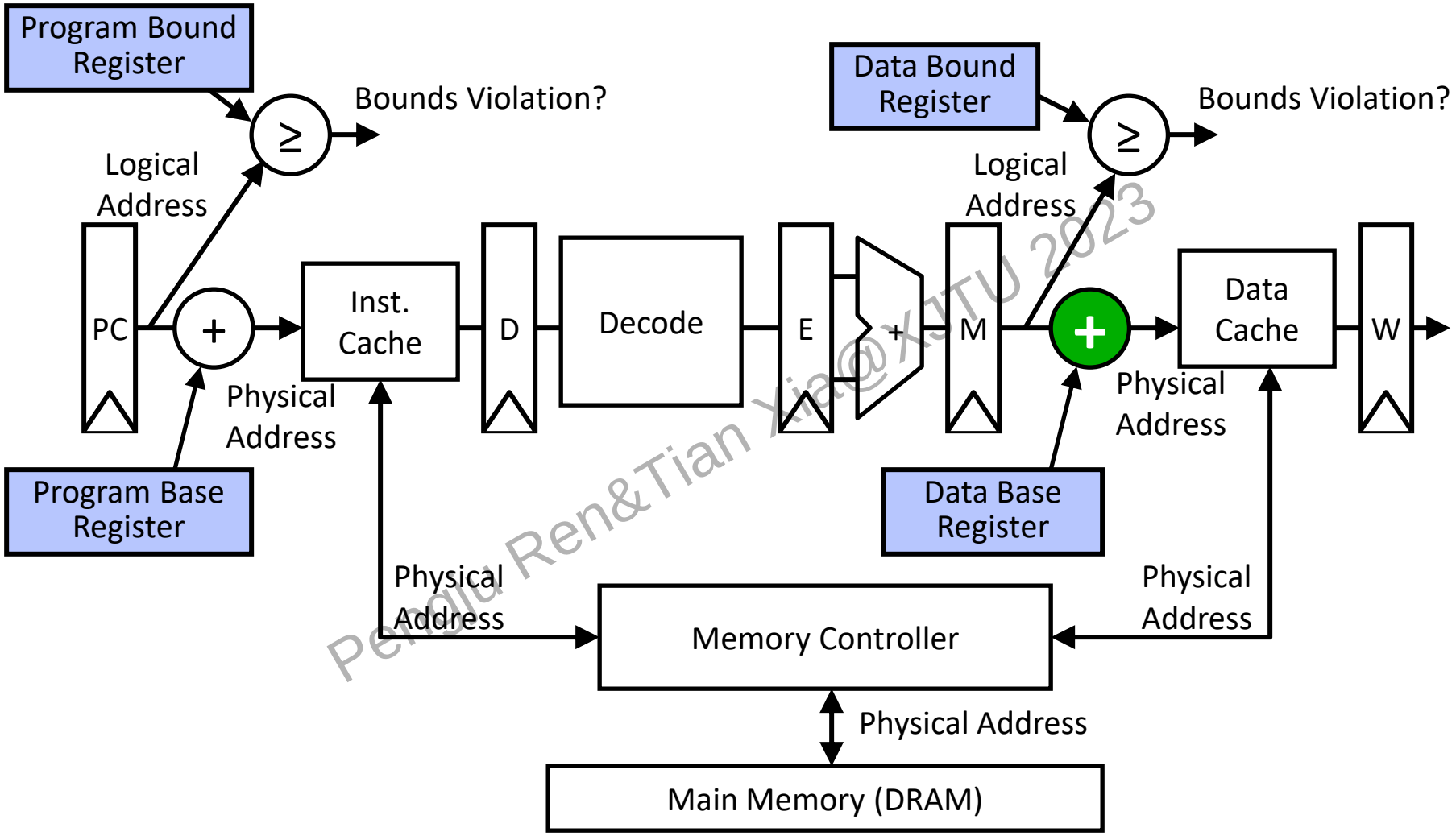


What is an advantage of this separation?

- Shared program segment
- Very fast

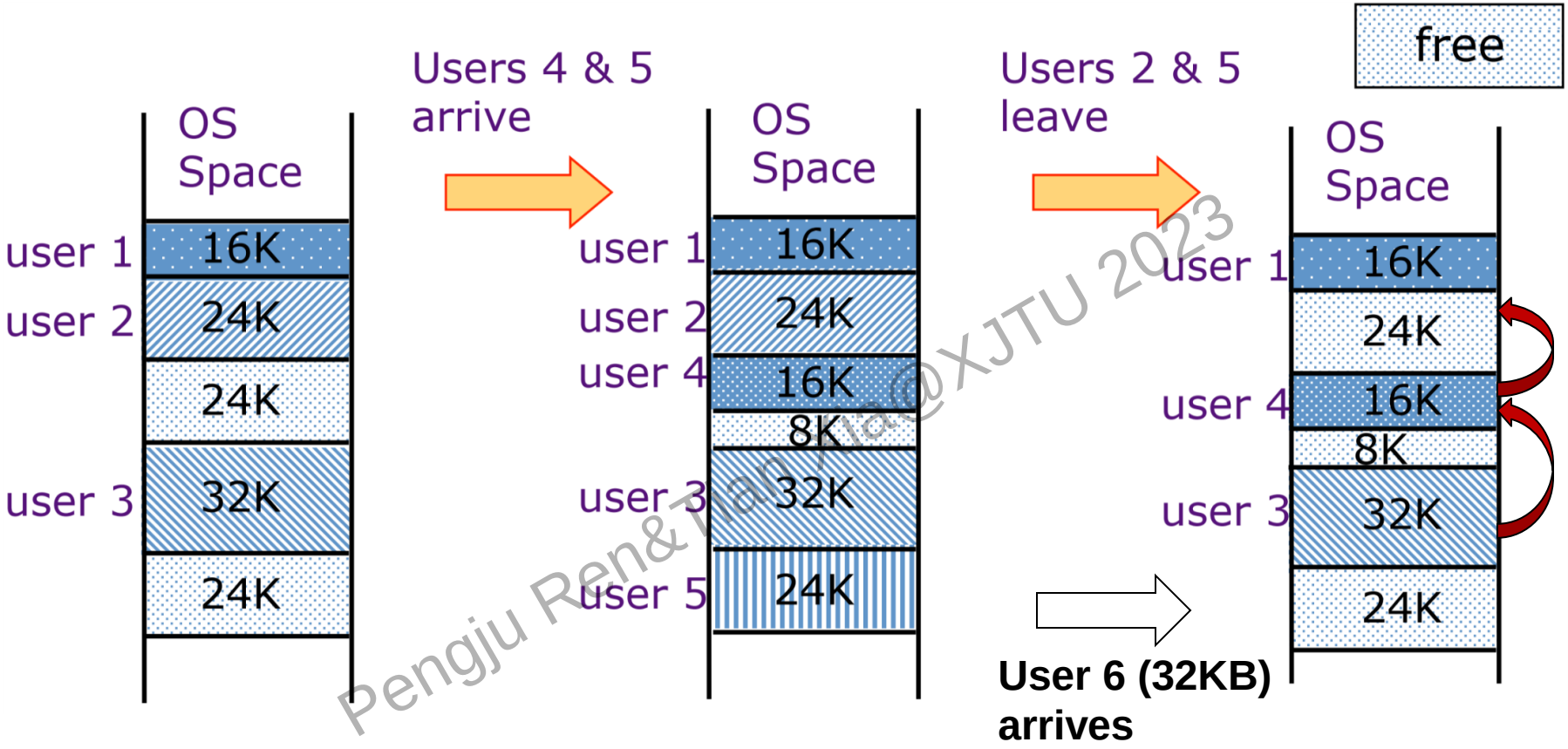
What about more base/bound pairs?

Base and Bound Machine



Can fold addition of base register into (register+immediate) address calculation using a carry-save adder (sums three numbers with only a few gate delays more than adding two numbers)

External Fragmentation with Segments

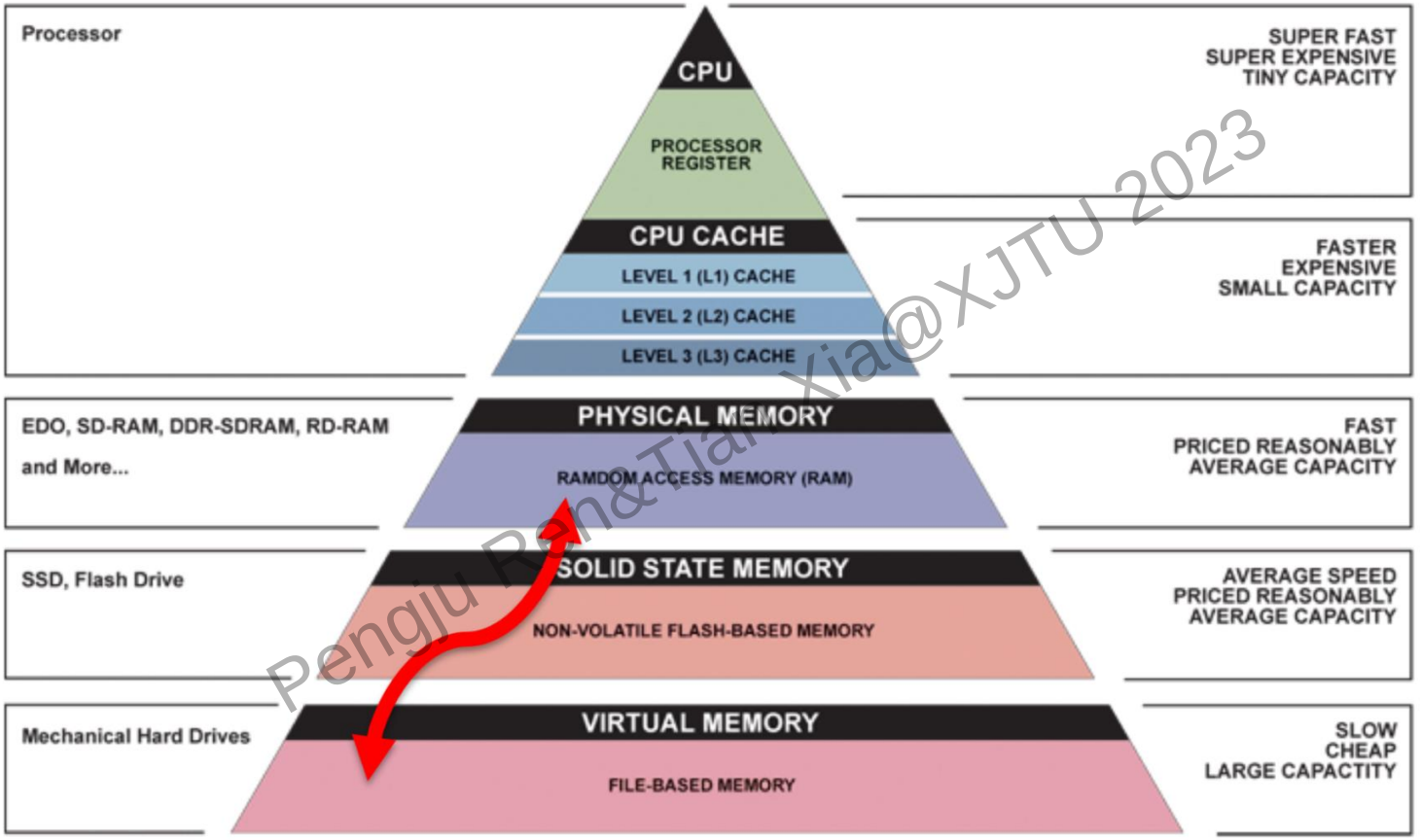


As users come and go, the storage is "fragmented"(external)

- **Plan ahead** to avoid bubbles
- Programs are **moved** around to compact the storage

Why do we need Virtual Address Space ?

Reason 1: Adding Disks to Hierarchy



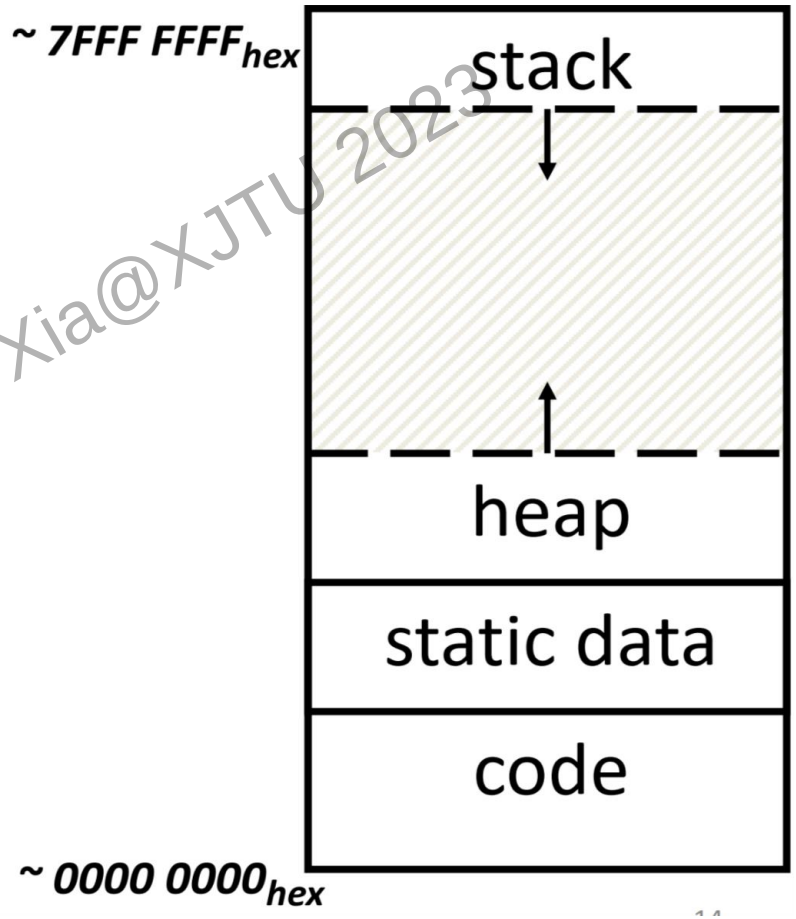
▲ Simplified Computer Memory Hierarchy
Illustration: Ryan J. Leng

Need to devise a mechanism to "**Connect**" memory and disk in the memory hierarchy.

Why do we need Virtual Address Space ?

Reason 2: Simplifying Memory for Apps

- Applications should see the straightforward memory layout we saw earlier ->
- User-space applications should think they own all of memory
- So we give them a **virtual** view of memory



Why do we need Virtual Address Space ?

Reason 3: Protection Between Processes

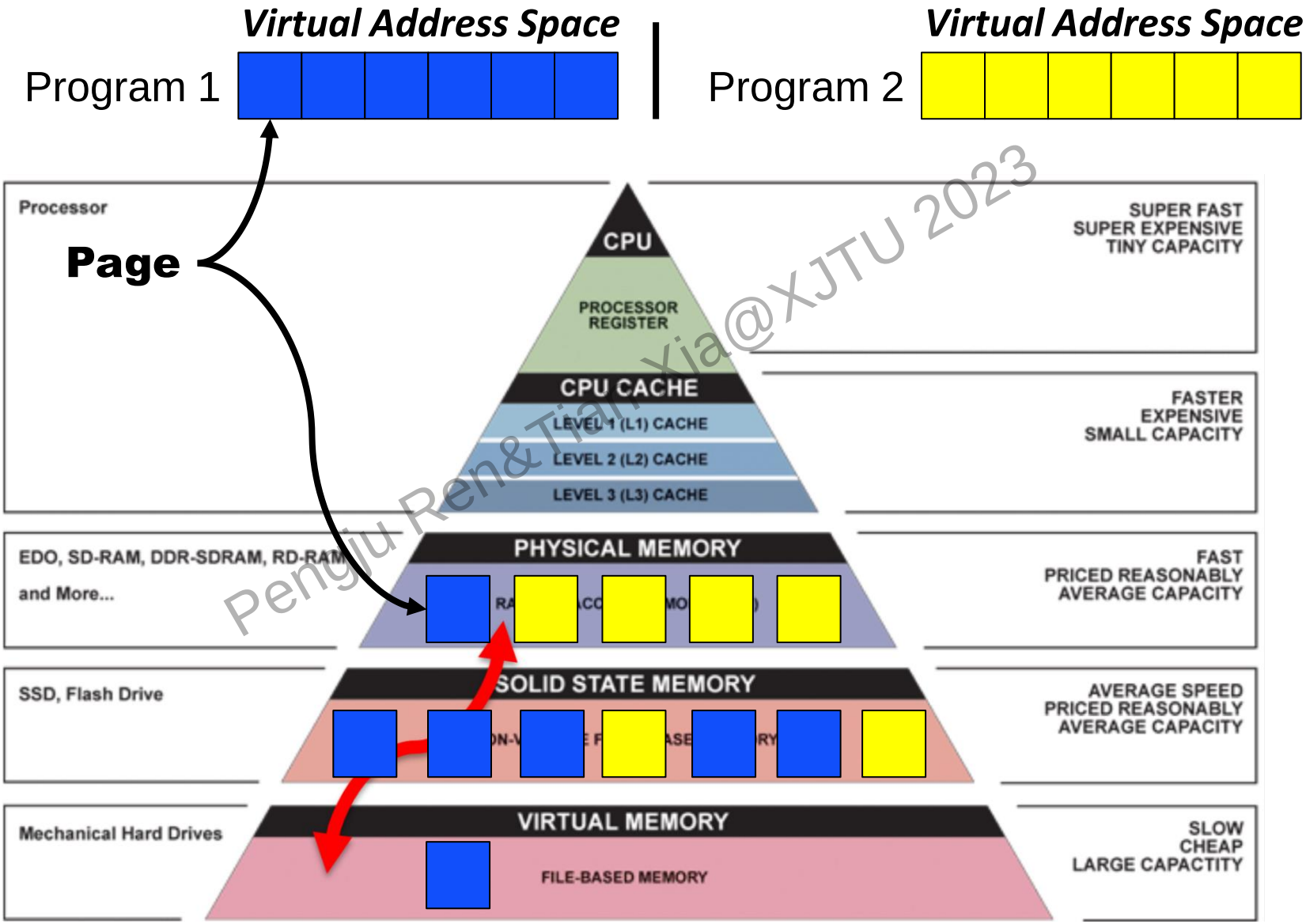
- With a bare system, addresses issued with loads/stores are real **physical** addresses
- This means any program can issue any address, therefore can access **any part of memory**, even areas which it doesn't own
  – Ex: The OS data structures
- We should send all addresses through a mechanism that the **OS controls**, before they make it out to DRAM
  - ***a translation and protection mechanism***

2 Parts to Modern Virtual Memory

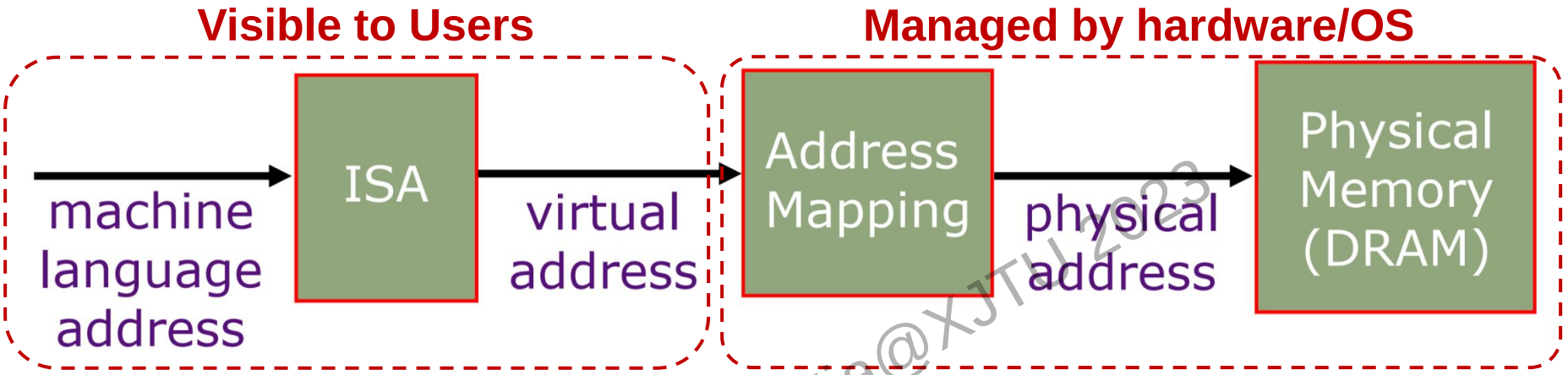
In a multi-tasking system, virtual memory supports the **illusion** of a **large**, **private**, and **uniform** memory space to each process

- Ingredient A: naming and protection
  - – each process sees a large, contiguous address space without holes ***(for convenience)***
  - – each process's memory is private, i.e., protected from access by other processes ***(for sharing and protection, Readable? Writeable? Executable?)***
- Ingredient B: demand paging ***(for hierarchy and efficiency)***
  - – location-independent programs
  - – capacity of secondary storage (swap space on disk)
  - – speed of primary storage (DRAM)

Why do we need Virtual Address Space ?



Names for Memory Locations



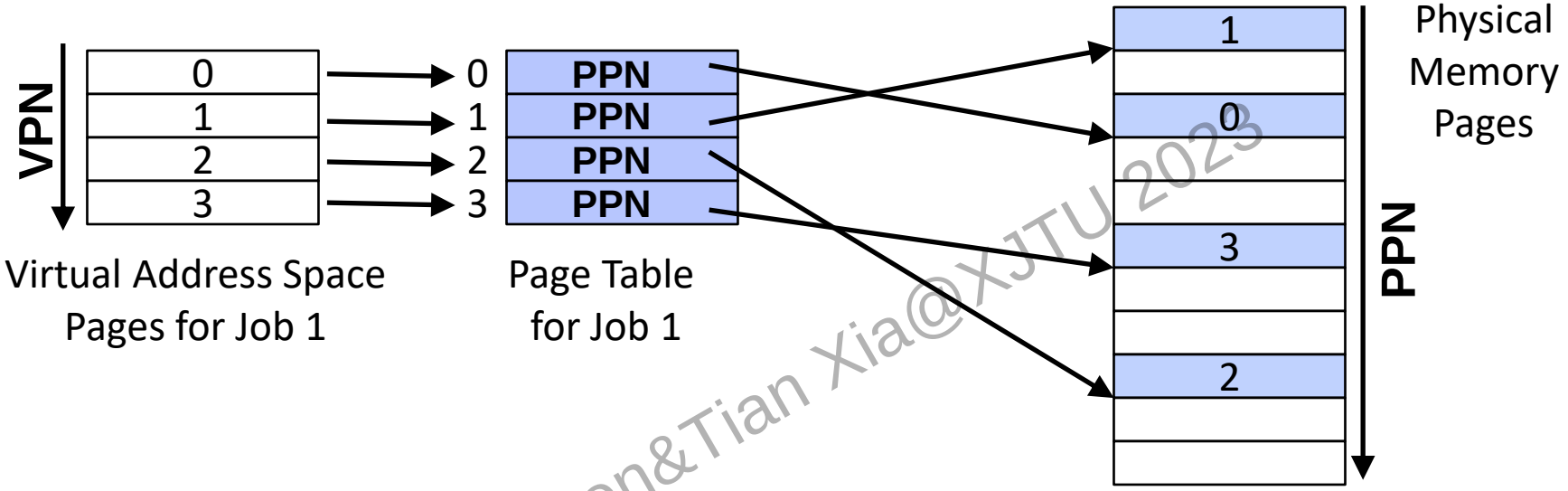
- Machine Language address
  - As specified in machine code
- Virtual Address (VA)
  - ISA specifies translation of machine code address into virtual address of program variable (sometime called effective address)
- Physical Address (PA)
  - Operating System specifies mapping of virtual address into name for a physical memory location

The Common Denominator : Address Translation

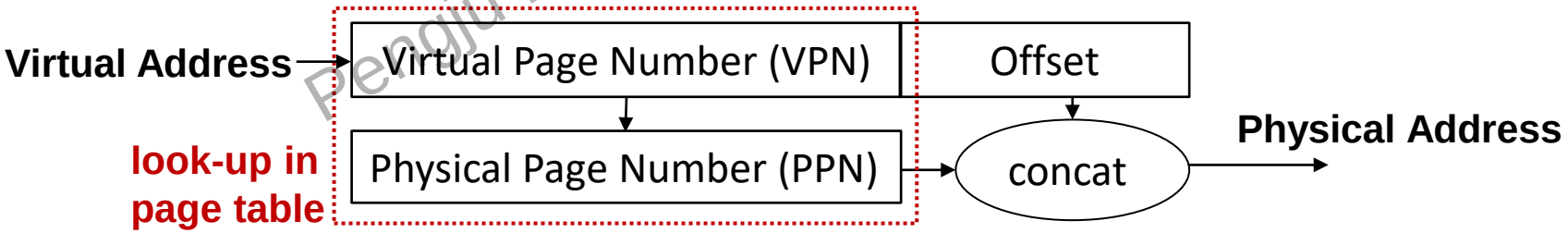
- Large, private, and uniform abstraction achieved through address translation
  - user process operates on virtual address (**VA**)
  - HW translates **VA** to physical address (**PA**) on every **memory reference**
- Through address translation
  - control which **physical locations** (DRAM and/or swap disk) can be referred to by a process
  - allow **dynamic allocation** and relocation of physical backing store (where in DRAM and/or swap disk)
- Address translation HW and policies **controlled by the OS** and protected from user

Paged Memory Systems (How)

- **Fixed-sized pages (mostly 4KB)** in virtual address space are mapped to physical address using **Page Table (PT)**

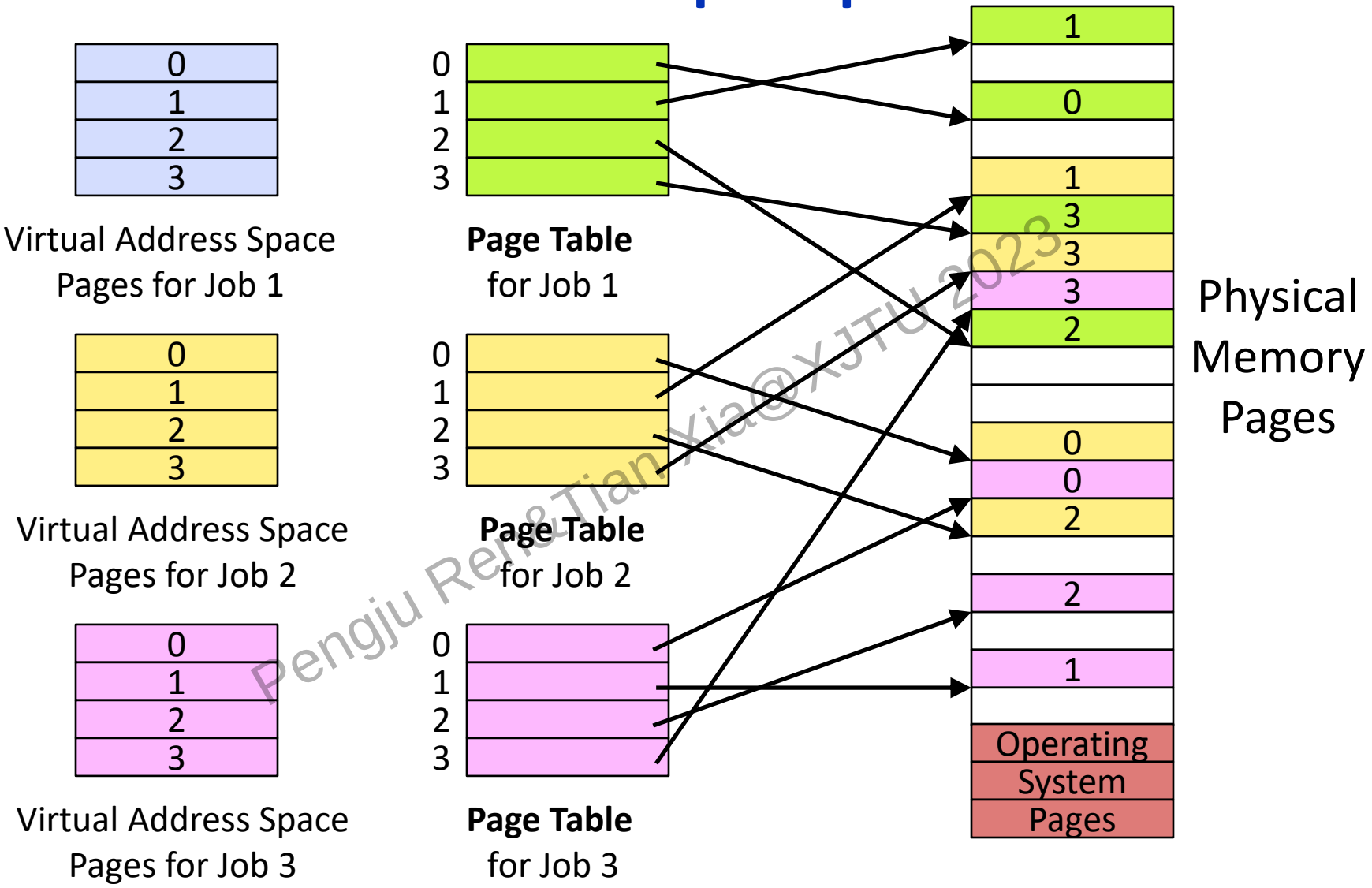


- Program-generated virtual address is split into **{VPN + offset}**



- For each program, an **independent** Page Table is maintained
- Paging makes it possible to store a large **contiguous** virtual memory space using **non-contiguous** physical memory pages

Private Address Space per User



- Each user has a **Page Table** contains an entry for each user page
- **Persistent OS** residing in memory to control all page tables

Paging Simplifies Memory Allocation

- Fixed-size pages can be kept on **OS free list** and allocated as needed to any process
- Process memory usage can easily **grow and shrink** dynamically
- Paging suffers from **internal fragmentation** *(inside Page)* where not all bytes on a page are used
  - Much less of an issue than external fragmentation or compaction for common page sizes (4-8KB)
  - But one reason that many oppose move to larger page sizes

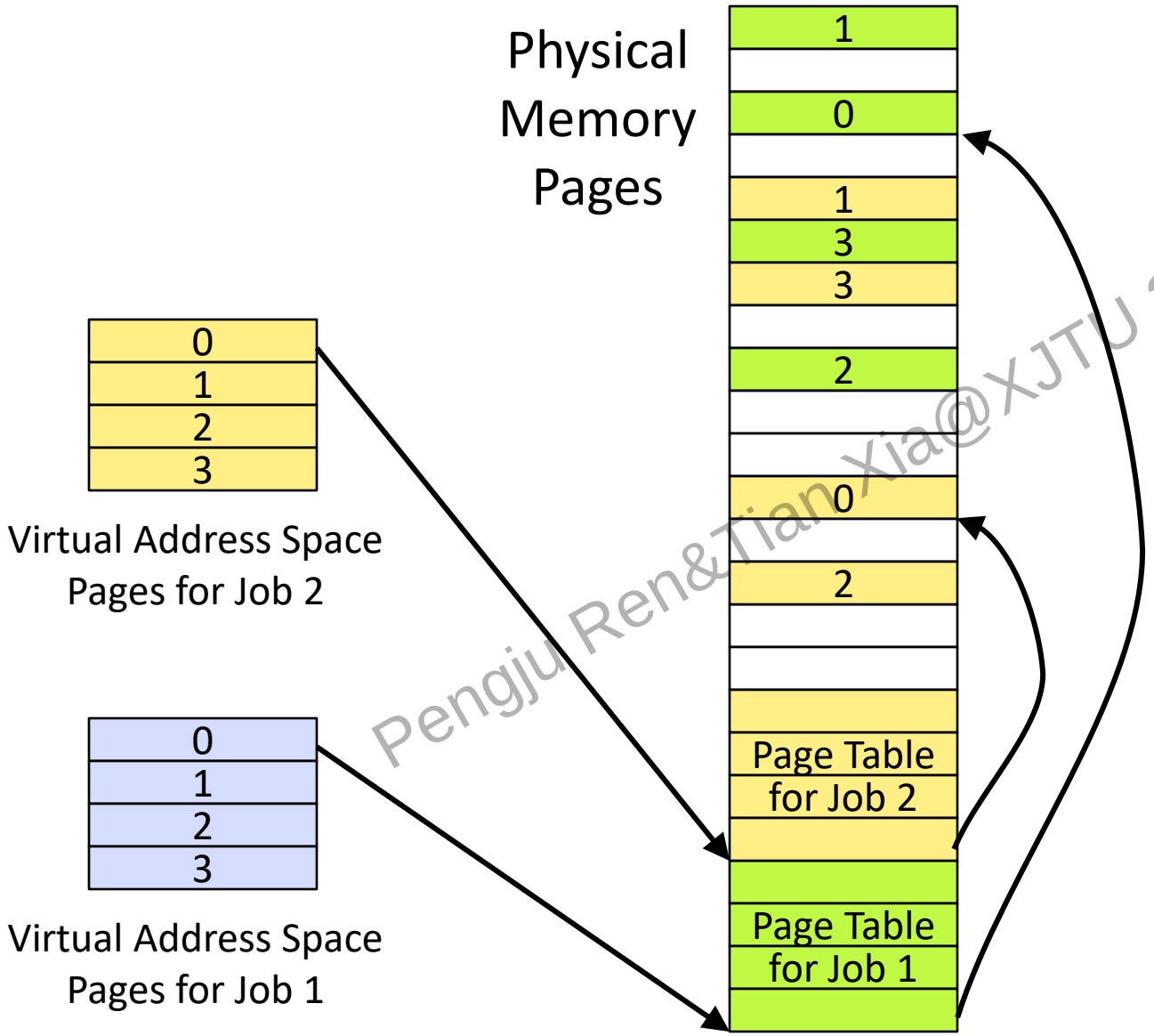
Coping with Limited Primary Storage

- Paging reduces fragmentation, but still many problems would not fit into primary memory, have to move data to/from secondary storage (drum, disk)
- Early approach:
  - **Manual overlays**, programmer explicitly copies code and data in and out of primary memory
    - Tedious coding, error-prone (jumping to non-resident code?)
    - IBM **Cell** microprocessor using in Playstation-3 had explicitly managed local store!
    - Many new "**deep learning**" accelerators have similar arch.
- Using virtual memory pages:
  - Put a physical page in **primary or secondary** storage wherever suitable
  - Maintain its position in a virtual memory **page table entry**.

Where should Page Tables Reside ?

- Space required by the page table (PT) is proportional to the address space, number of users, …
  - Space requirement is large (**4GB** space = **1M** PT Entries of 4KB pages for **each user** = **4MB** for 32-bit entry)
- Bad Idea: Keep PT of current user in special registers
  - May not be feasible for large page tables
  - Too expensive to keep in registers
  - Increases the cost of context swap
- Good Idea: Keep PTs in the main memory
  - Use one **Page Table Base Register** (**PTBR**) to hold PT's location in the main memory
  - Needs **one additional reference** to retrieve the page base address and another to access the data word
  - **Double the number of memory references!**

Page Tables Live in Memory

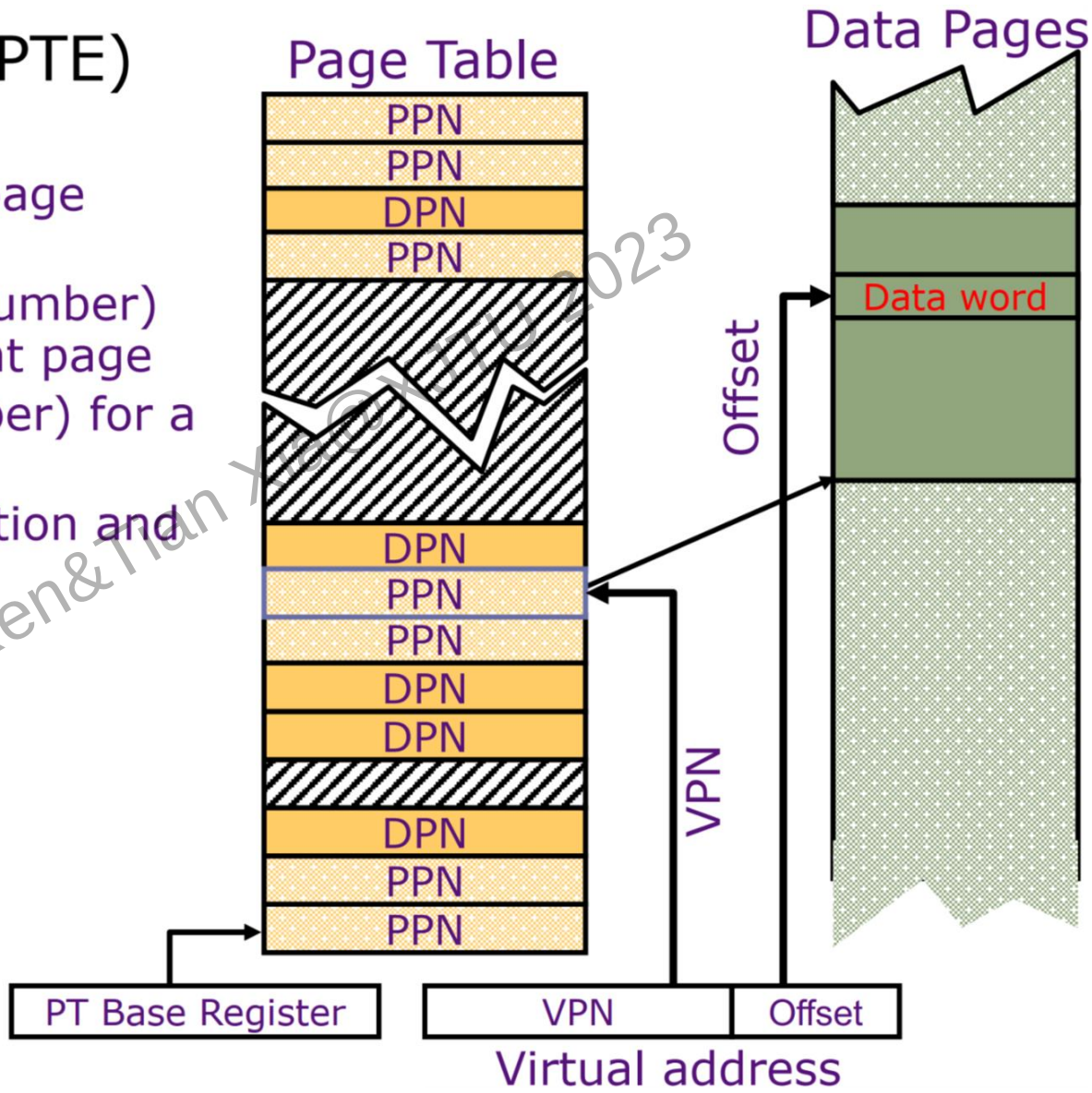


*Simple linear page tables are too large, so **hierarchical page tables** are commonly used (see later)*

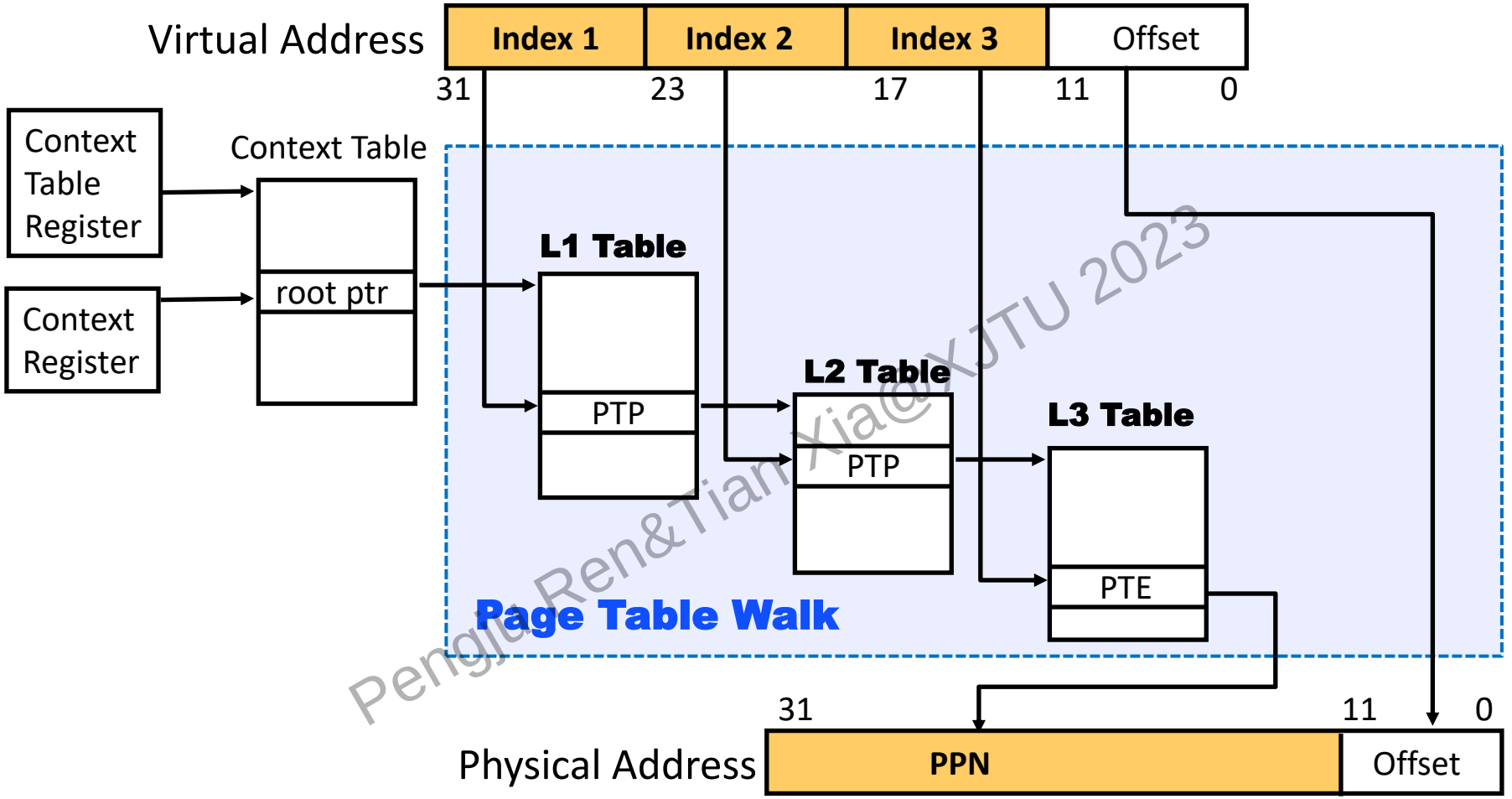
*Common for modern OS to place page tables in kernel's **virtual memory** (page tables can be swapped to secondary storage)*

Linear Page Table

- Page Table Entry (PTE) contains:
  - A bit to indicate if a page exists
  - PPN (physical page number) for a memory-resident page
  - DPN (disk page number) for a page on the disk
  - Status bits for protection and usage

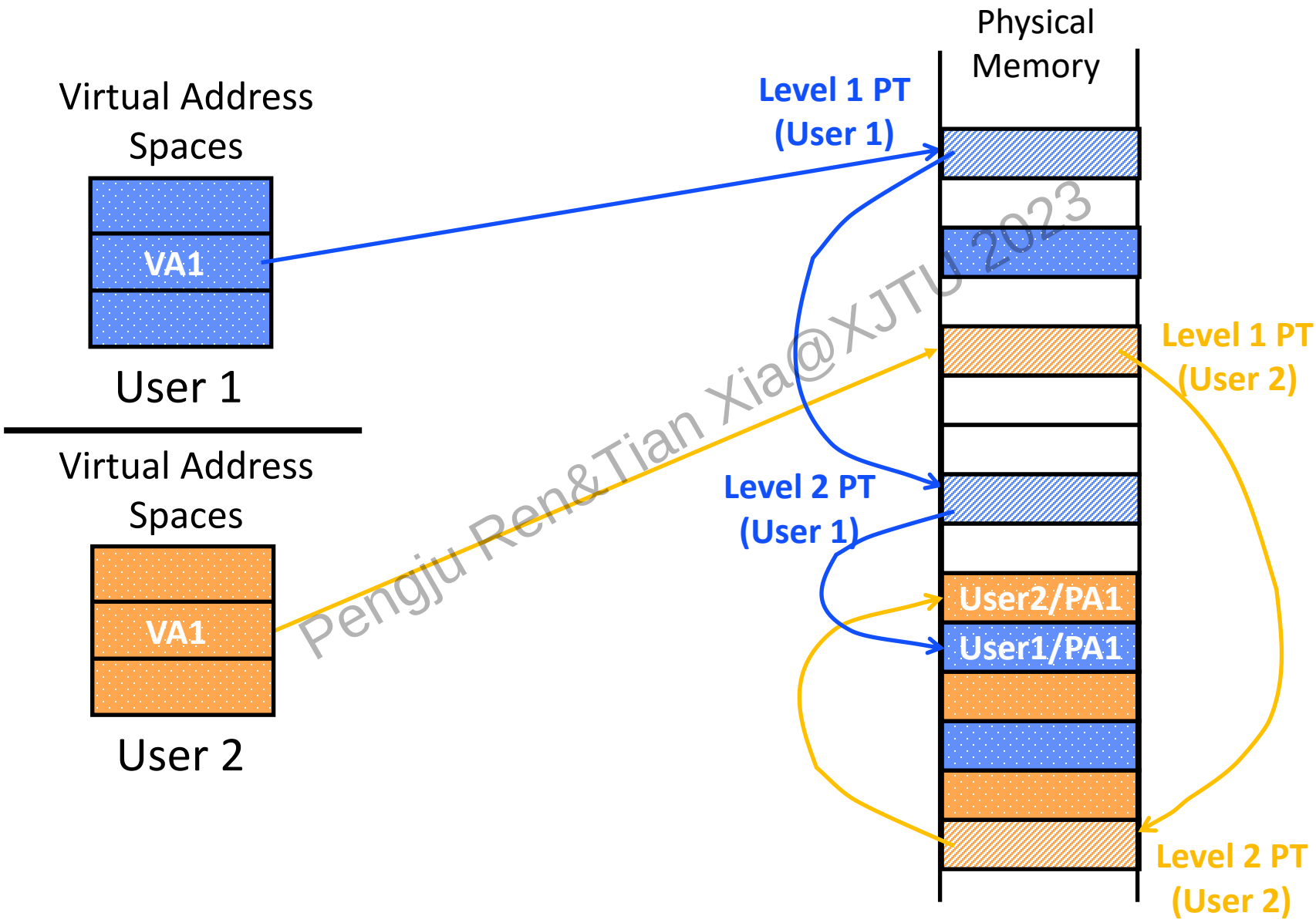


Hierarchical Page Table Walk: SPARC v8



- Hierarchical Page Table is a **tree** structure, PTBR is the **root**.
- Only create page table when **necessary**, reduces memory footprint
- Termed as **page table walk** , usually performed in **hardware unit.**

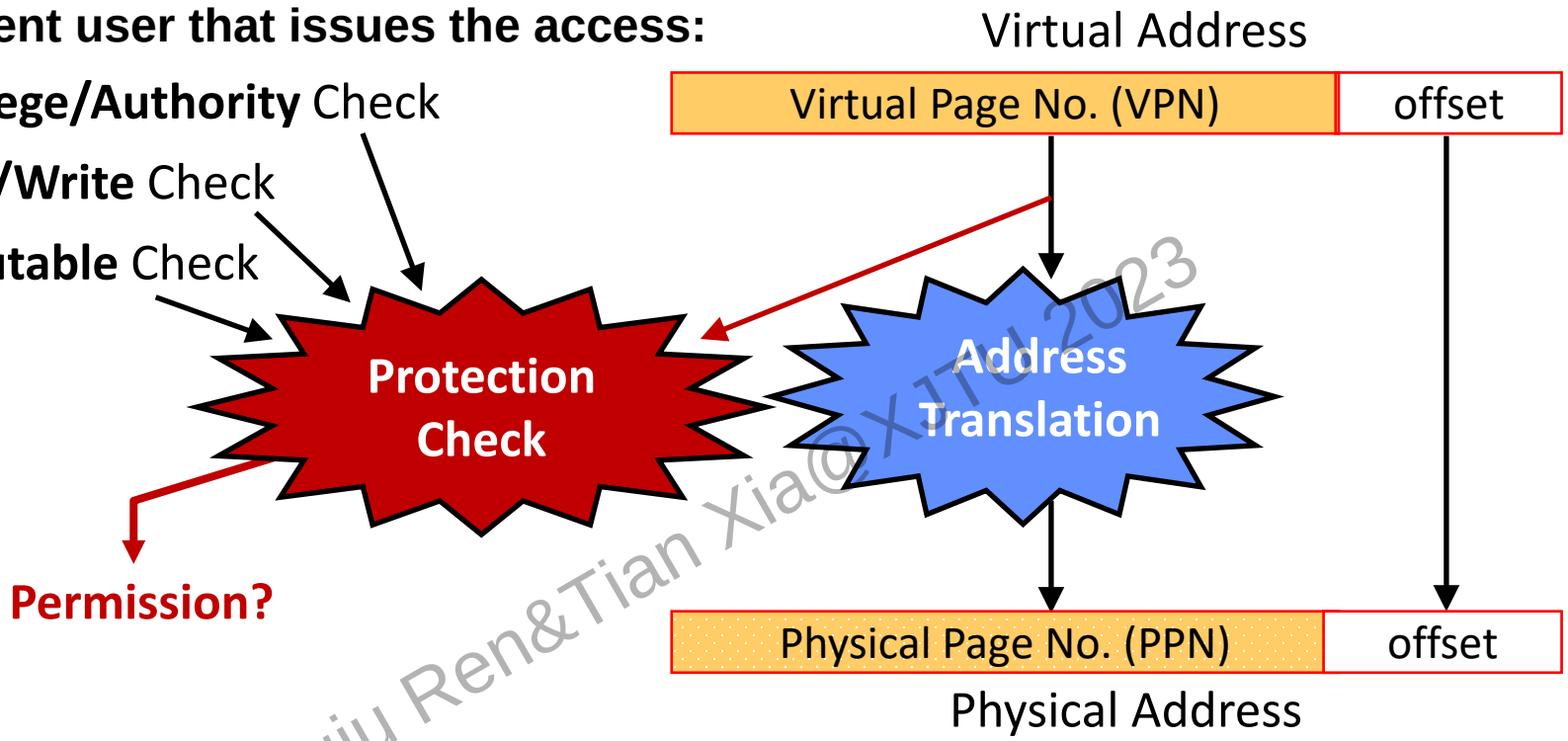
Two-Level Page Tables in Physical Memory



Address Translation & Protection Check

For current user that issues the access:

- **Privilege/Authority** Check
- **Read/Write** Check
- **Executable** Check

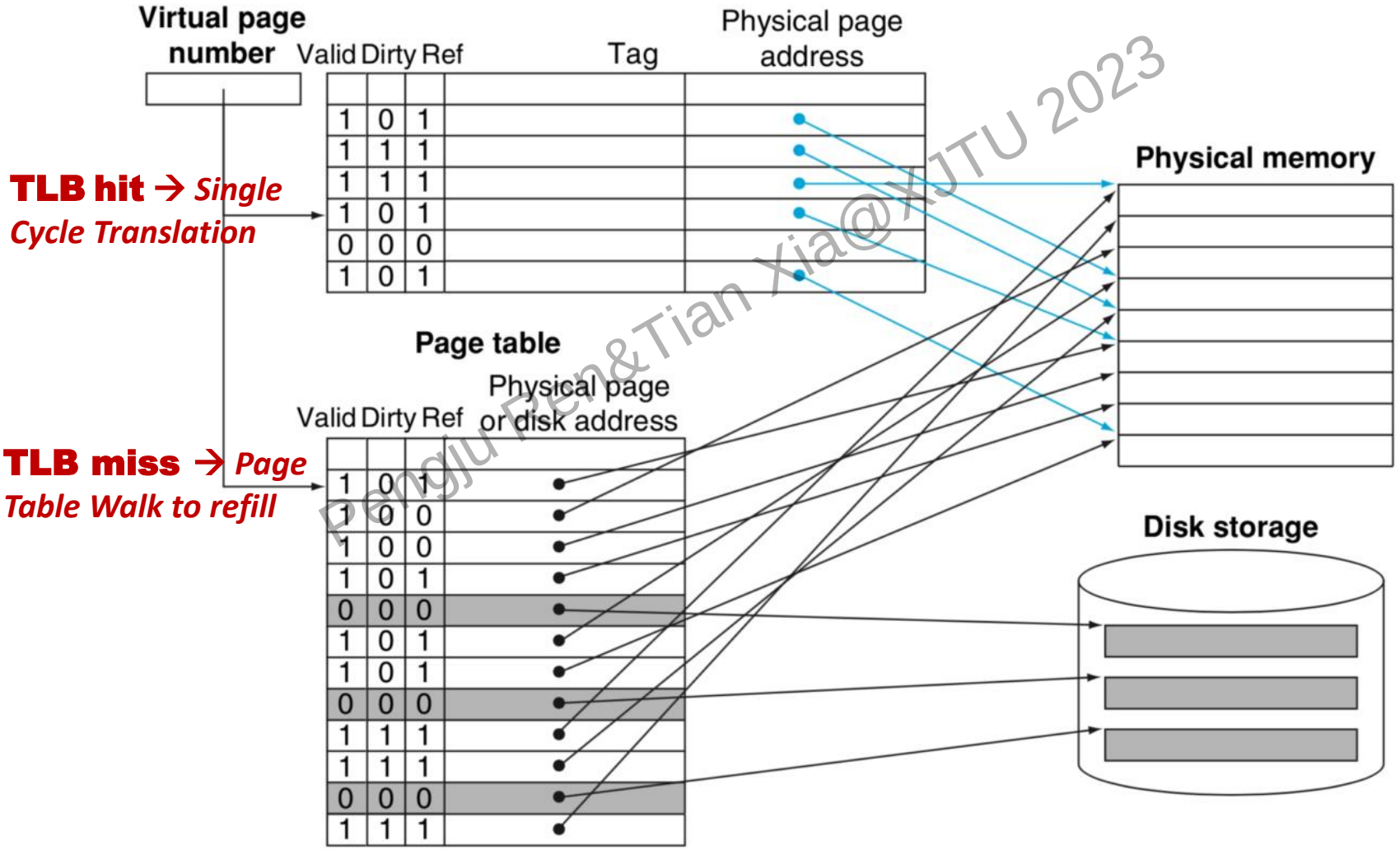


- Every instruction/data access needs address translation and **protection checks**
- Address translation is very **expensive**!
  In a two-level page table, each reference requires several memory accesses

Translation-Lookaside Buffers (TLB)

*A good VM design needs to be **fast** (~ one cycle) and space **efficient***

Idea: ***Cache** the address translation of **frequently used** pages*

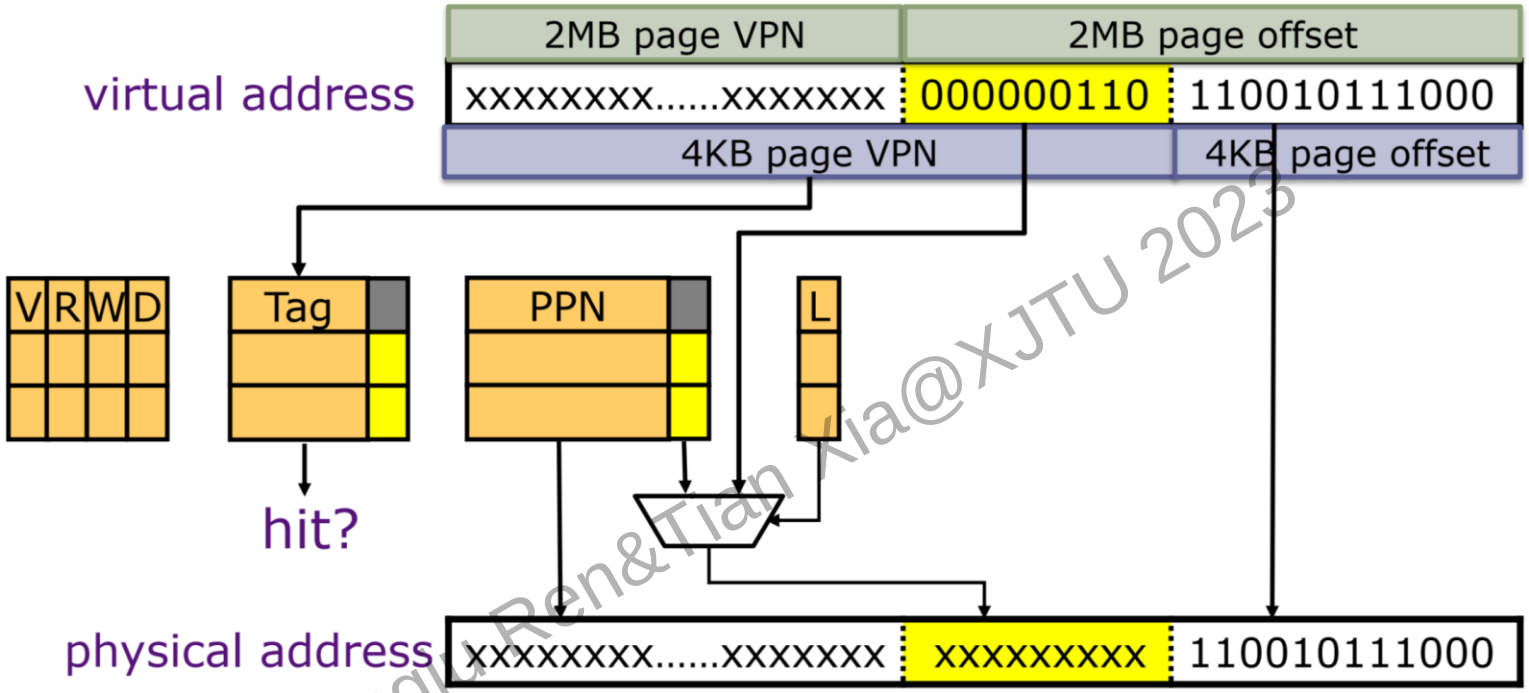


TLB Designs

- Typically **32-128** entries, usually **fully associative**
  - Each entry maps a large page, hence less spatial locality across pages ➔ more likely that two entries conflict
  - Sometimes larger TLBs (256-512 entries) are 4-8 way set-associative
  - Larger systems sometimes have multi-level (L1 and L2) TLBs
- **Replacement policy**: Random or FIFO
- **TLB Reach**: Size of largest virtual address space that can be simultaneously mapped by TLB
  - Example: 64 TLB entries, 4KB pages, one page per entry
  - TLB Reach = ***64 entries * 4 KB = 256 KB (if contiguous)***?

Variable-Size Page TLB

Some Systems support multiple page sizes



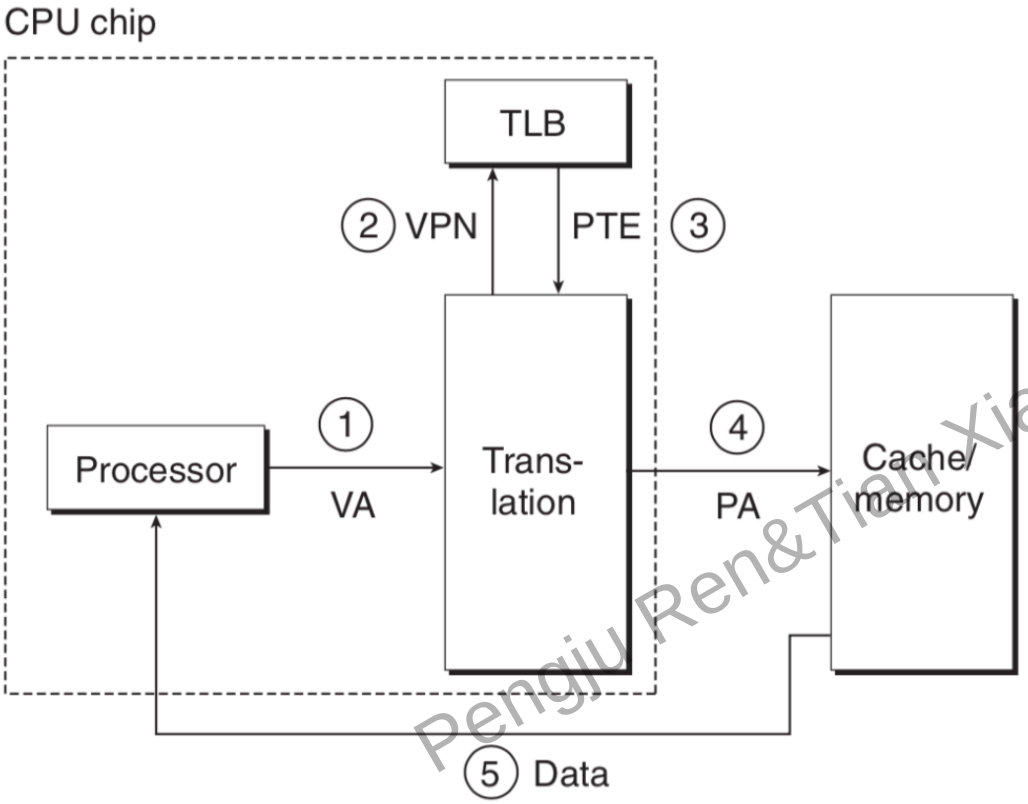
Page mask	Page Size	Pagemask	Page Size
0_0000_0000	4KB	0_0011_1111	256KB
0_0000_0011	16KB	0_1111_1111	1MB
0_0000_1111	64KB	1_1111_1111	2MB

MIPS using Pagemask mark different Page size (OS manage)

Handling a TLB Miss

- TLB misses when:
  - The page exists in memory → just add the missing entry in TLB.
  - The page doesn't exist in memory→ transfer control to the OS
- Software (*MIPS, Alpha*)
  - TLB miss causes an **TLB miss exception** and the operating system walks the page tables and reloads TLB.
  - **Very expensive** on out-of-order superscalar processor as requires a flush of pipeline to jump to trap handler.
- Hardware (*SPARC v8, x86, PowerPC, RISC-V*)
  - A **memory management unit (MMU)** walks the page tables and reloads the TLB.
  - If a missing (data or PT) page is encountered during the TLB reloading, MMU gives up and signals a **Page Fault exception** for the original instruction.

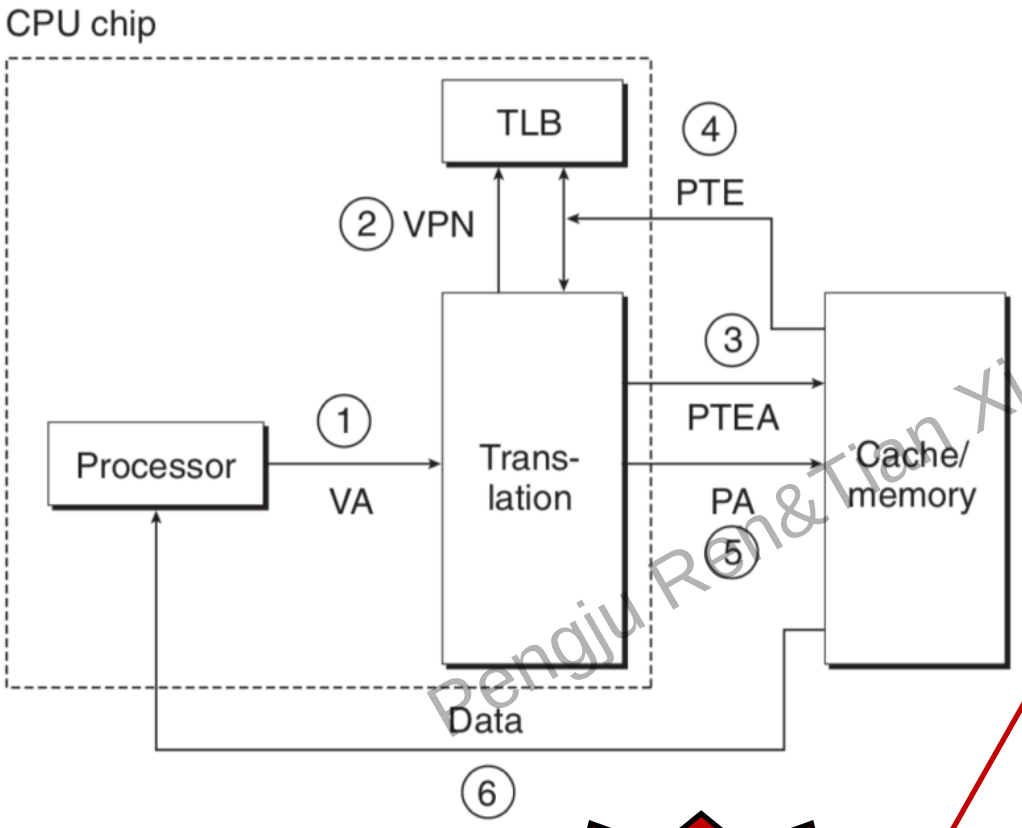
Handling a TLB Miss



TLB Hit

① Processor sends Virtual Address (VA)
② Extract Virtual Page Number(VPN) from VA. Query TLB using VPN.
③ TLB returns Page Table Entry (PTE).
④ Combine PTE with Page Offset to get Physical Address (PA). Query Cache using PA.
⑤ Send data to processor

Handling a TLB Miss



What if there is no valid PTE?

TLB Miss

① Processor sends Virtual Address (VA)

② Extract Virtual Page Number (VPN) from VA. Query TLB using VPN.

③ TLB miss. Query Memory using Page Table Entry Address (PTEA) to get Page Table Entry (PTE)

④ Save VPN to PTE mapping in TLB.

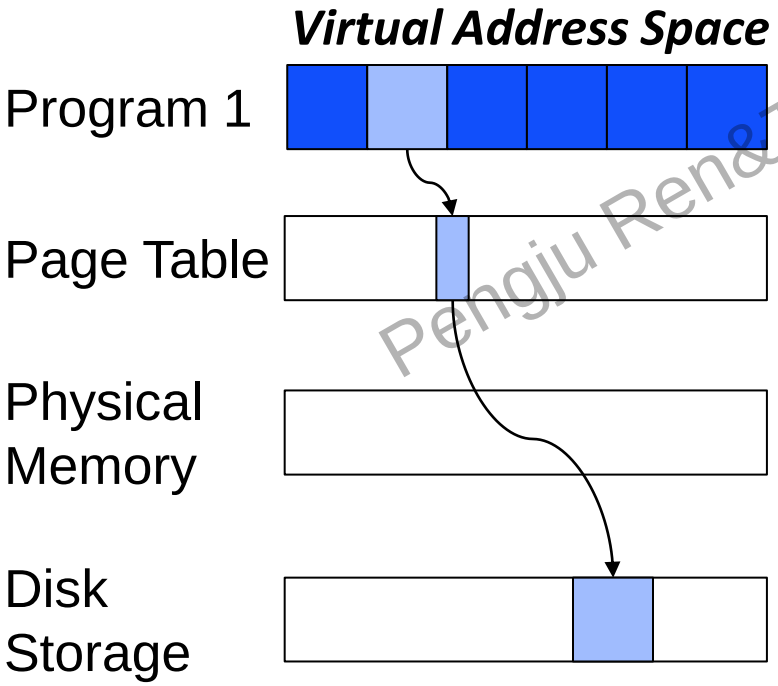
⑤ Combine PTE with Page Offset to get Physical Address (PA). Query Cache using PA.

⑥ Send data to processor.

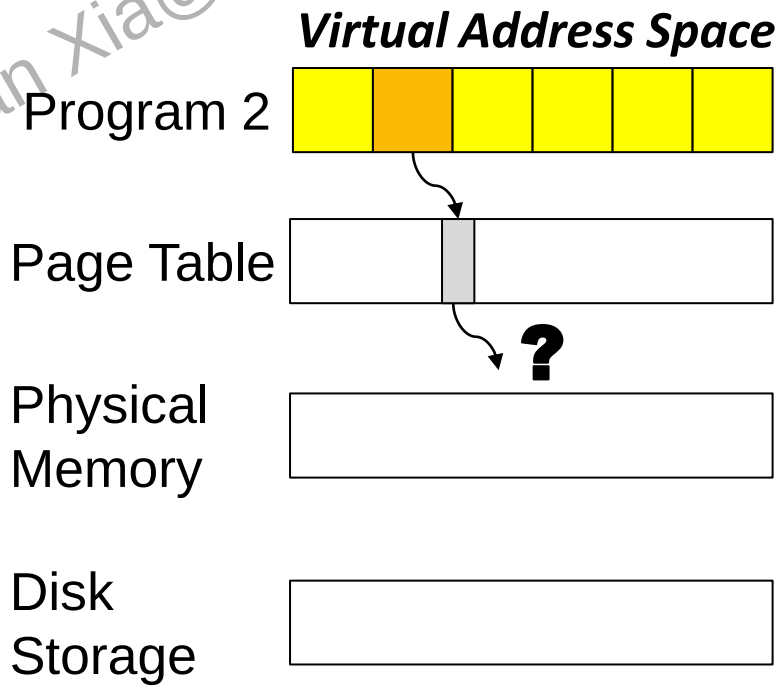
Page Fault Exception

- Occurs when an instruction references a memory page that is **not in main memory.**

Case 1: Page is swapped to secondary storage (e.g. disk)



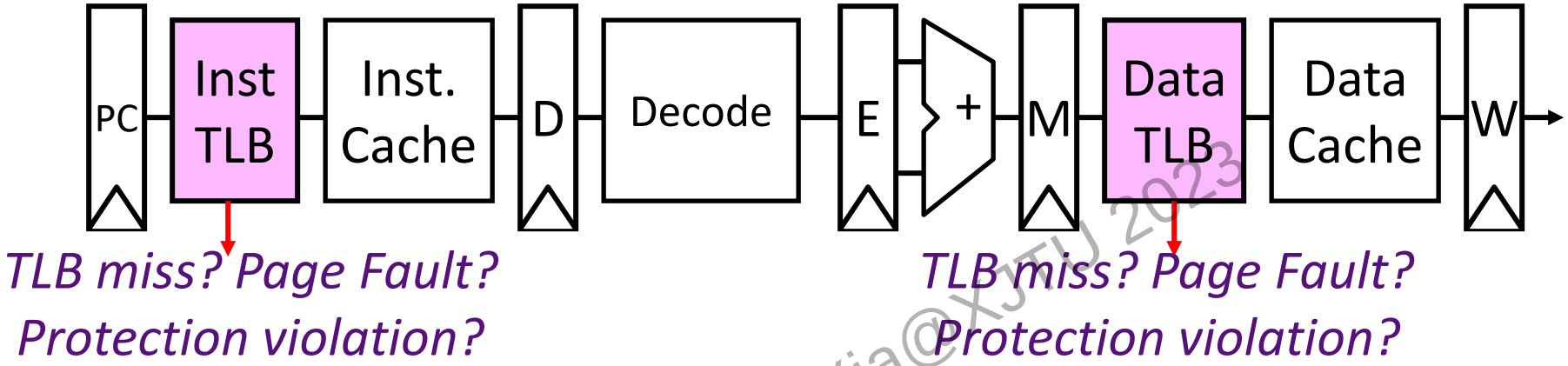
Case 2: Page is virtually allocated but not really created (e.g. malloc)



Page Fault Exception

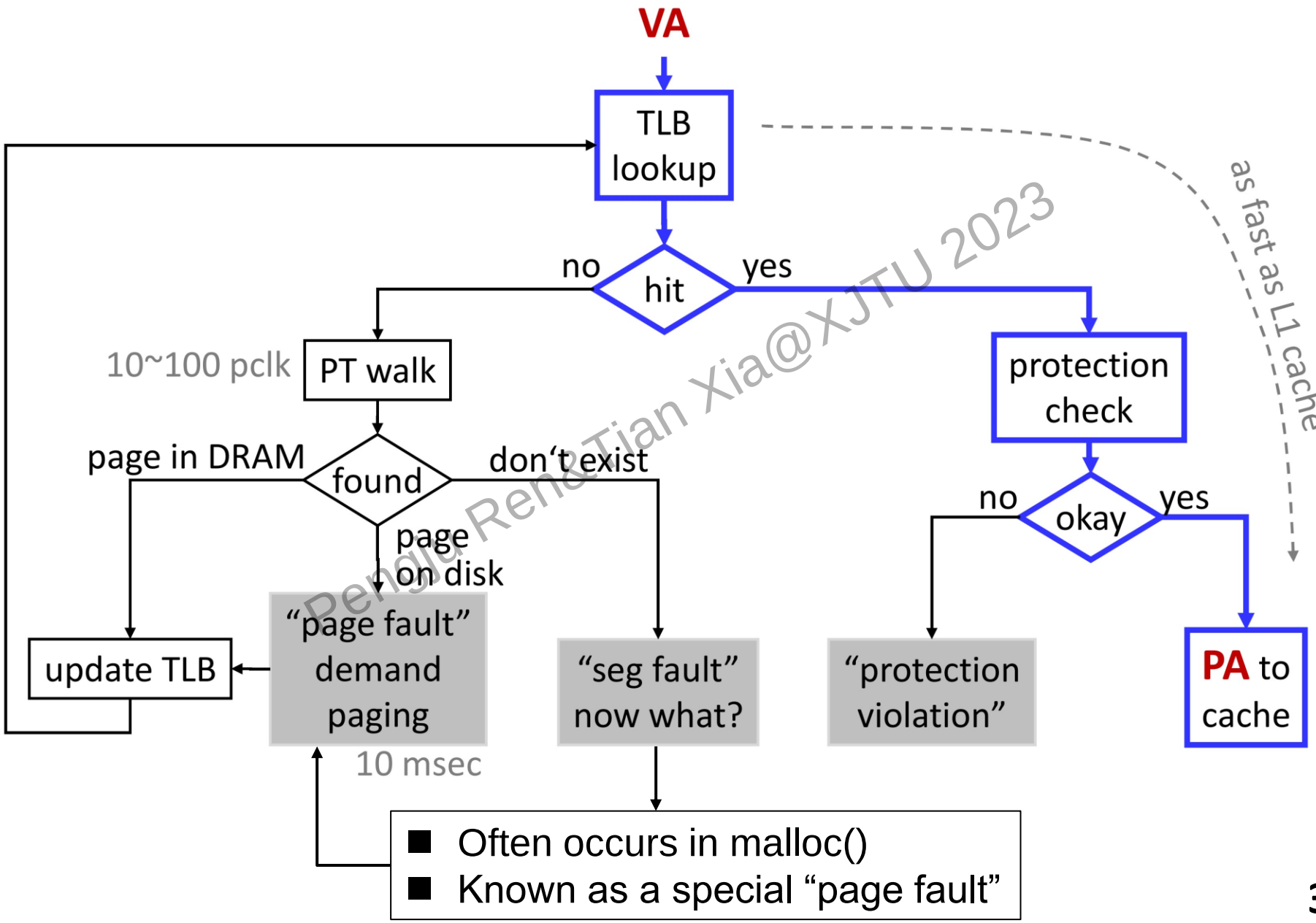
- Since it takes a long time to transfer a page (msecs), page faults are handled completely in software by OS
- **Page Fault Handler (of OS)** does the following:
  - Assign an unused page in DRAM
    - **If no unused page is left**, a page currently in DRAM is swapped out
      - Replacement policy: Pseudo-LRU, implemented in software
      - If the replaced page is 'dirty', write it back to
      - page table entry that maps that VPN->PPN is marked as invalid/DPN
  - If virtual page exist in disk, Initiate transfer of the requested page from disk to DRAM, assigning to an unused page
    - Another job may be run on the CPU while the first job waits for the requested page to be read from disk
  - Page table entry of the requested page is updated with a (now) valid PPN
  - Return and re-execute the exception-causing instruction
    - Need for precise exceptions

Handling VM-related exceptions

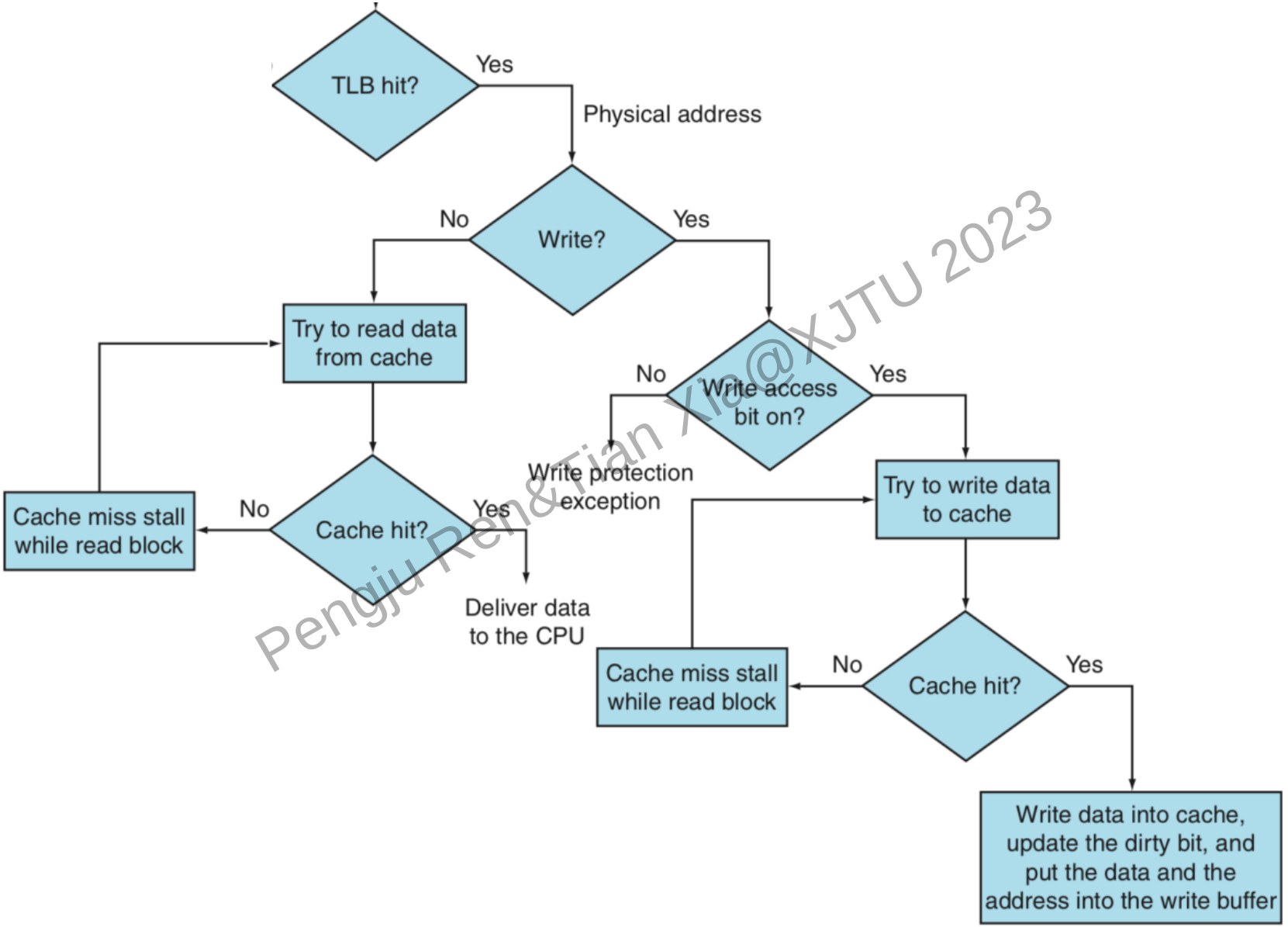


- Handling **TLB miss** needs a hardware or software mechanism to refill TLB（usually hardware as MMU）
- Handling **Page Fault** (e.g., page is on disk) needs *restartable* exception so software handler can resume after retrieving page
  - **Precise exceptions** are easy to restart
  - Can be imprecise but restartable, but this complicates OS software
- A **protection violation** may abort process
  - But often handled the same as a page fault

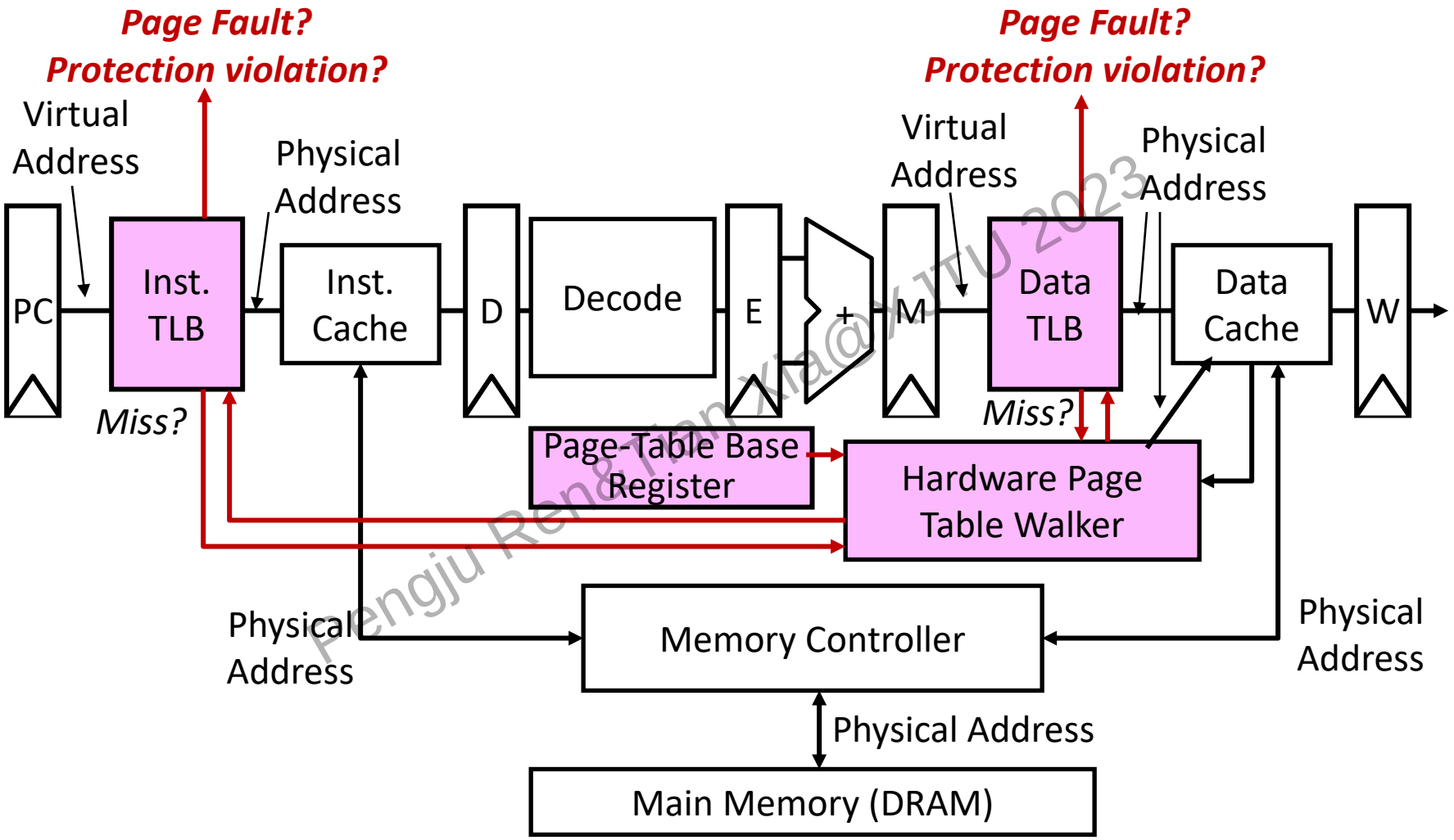
Address Translation – Putting it all together



Address Translation – Putting it all together

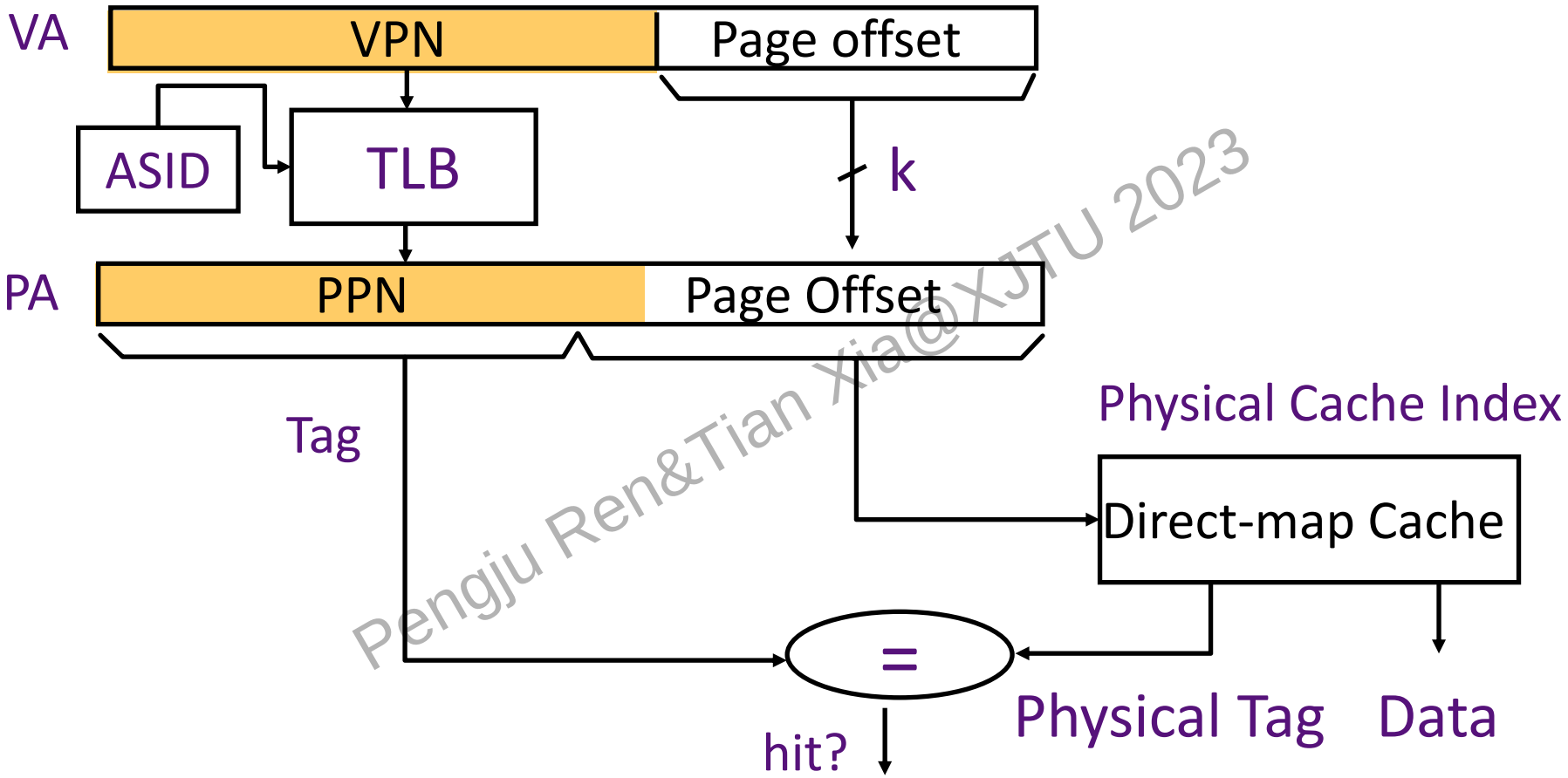


Page-Based Virtual-Memory Machine
(Hardware Page-Table Walk)



- MMU uses untranslated physical memory to access page tables

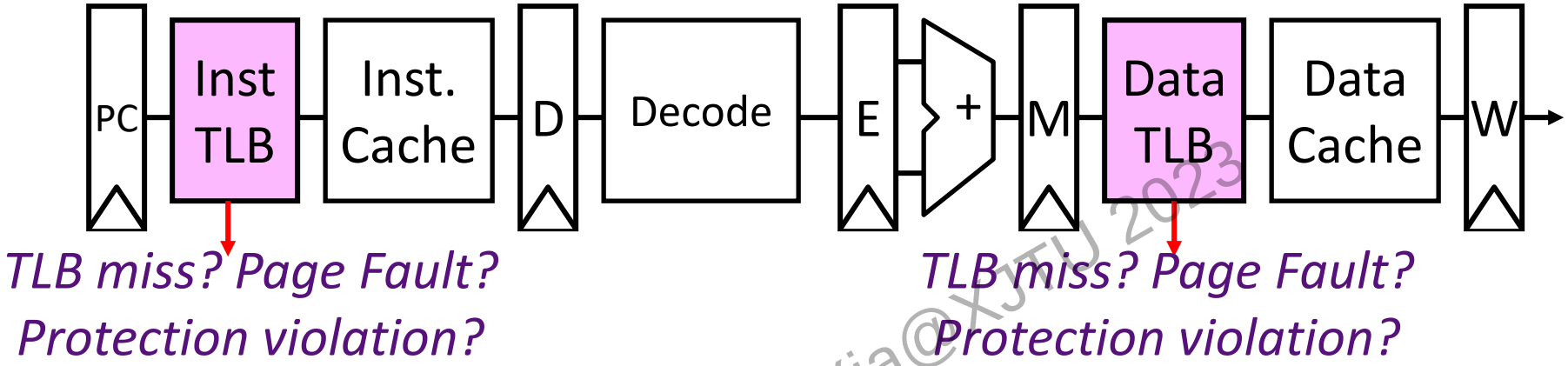
Sequential Access to TLB & Cache (Physical Index/Physical Tag, PIPT)



Adding one more stage for TLB access will increase:

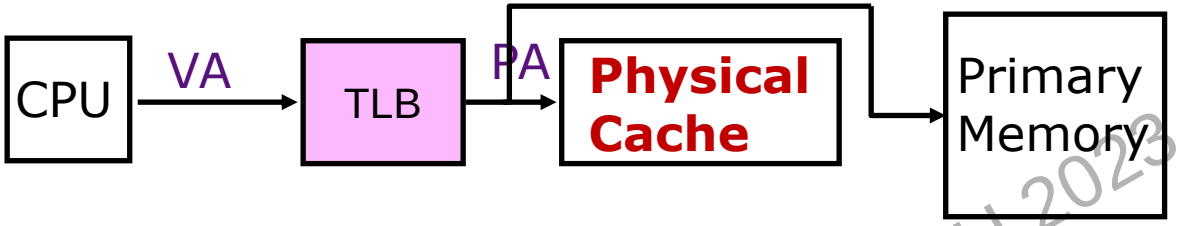
- For I-Cache the miss prediction penalty
- D-Cache load latency (**critical path!**)

Address Translation in CPU Pipeline

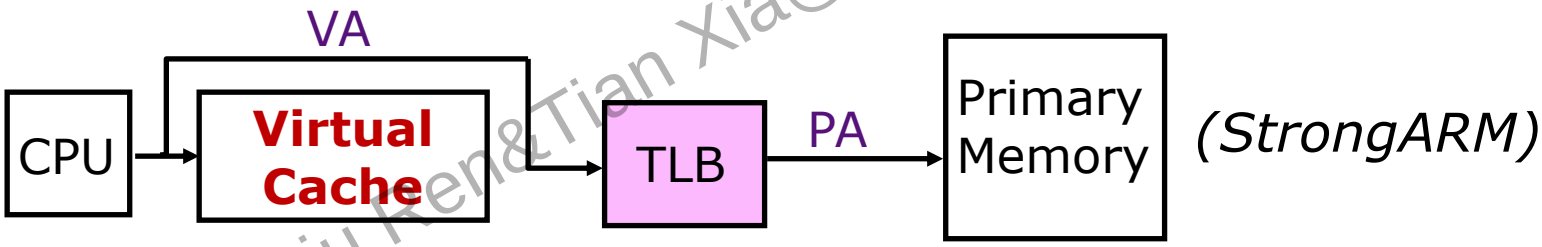


- Need to cope with additional latency of TLB:
  - slow down the clock?
    - Unacceptable for modern CPUs
  - pipeline the TLB and cache access?
    - Sub-optimal solution (still long latency)
  - virtual address caches
  - parallel TLB/cache access

Virtual-Address Caches



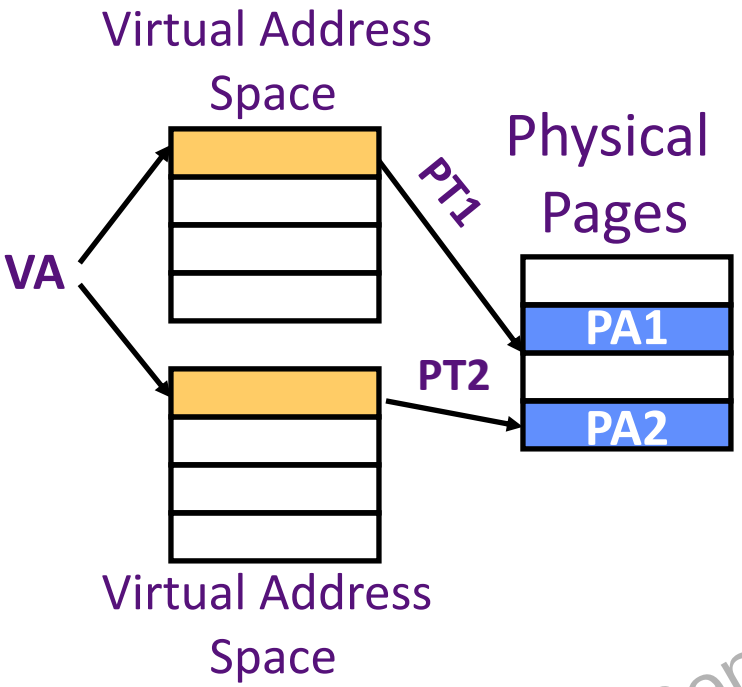
Alternative: place the cache before the TLB



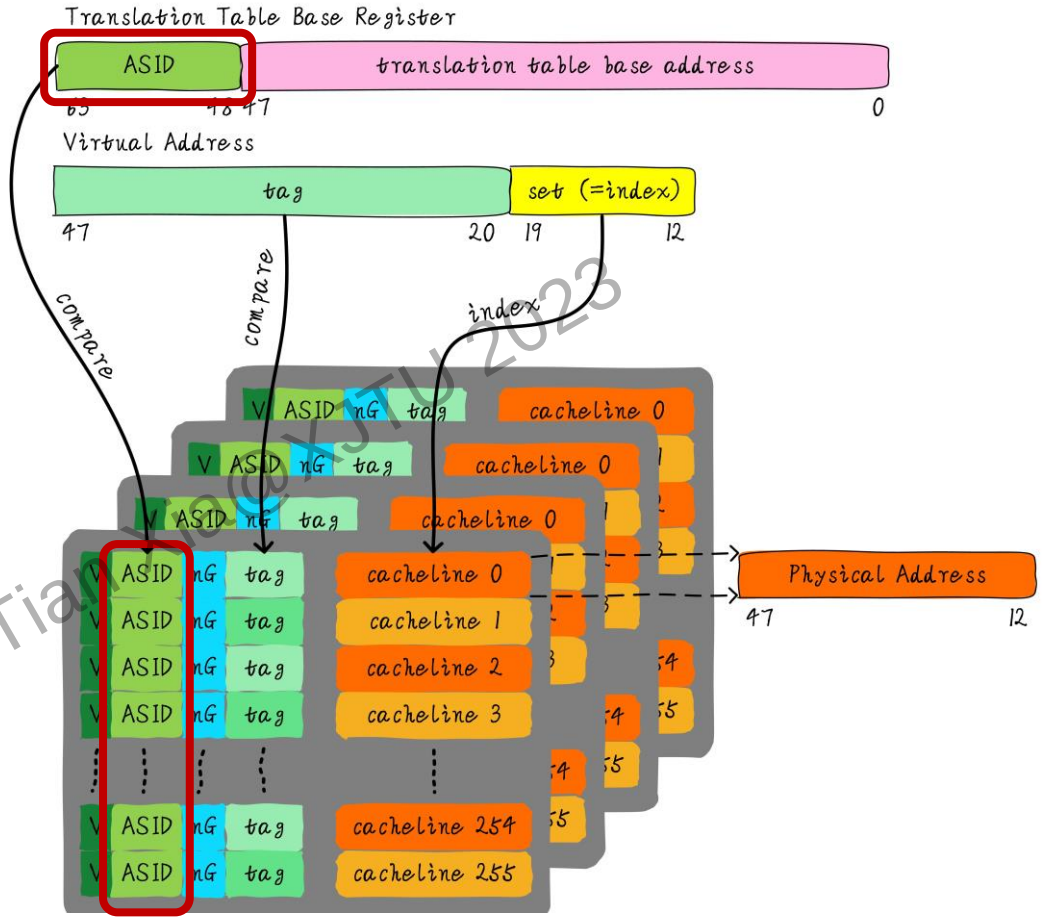
(StrongARM)

- one-step process in case of a hit (+)
- ***Homonym*** *and* ***aliasing*** *problems* (-)

Homonym in Virtual Address

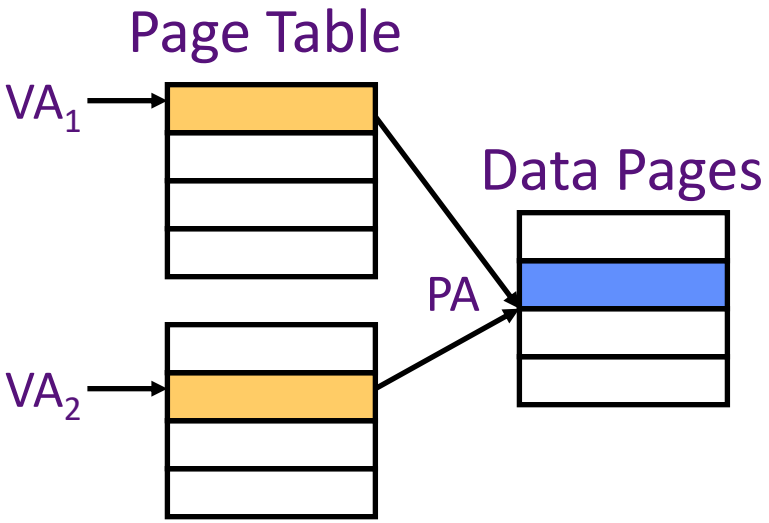


Same virtual address mapped to two physical pages

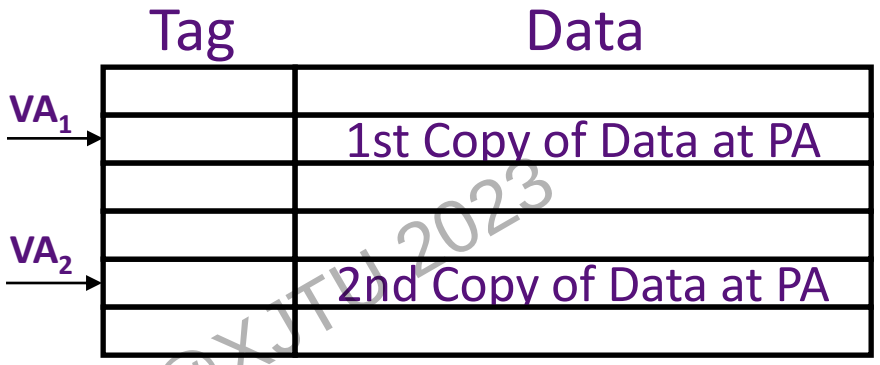


- Conflicting virtually-tagged entry (both **TLB** and VIVT **cache**)
- Software (OS): Cache and TLB needs to be flushed on a **context switch**
- Hardware: add Address Space Identifier (**ASID**) into Tags

Aliasing in Virtual-Address Caches



Two virtual pages share one physical page



Virtual cache can have two copies of same physical data. Writes to one copy not visible to reads of other!

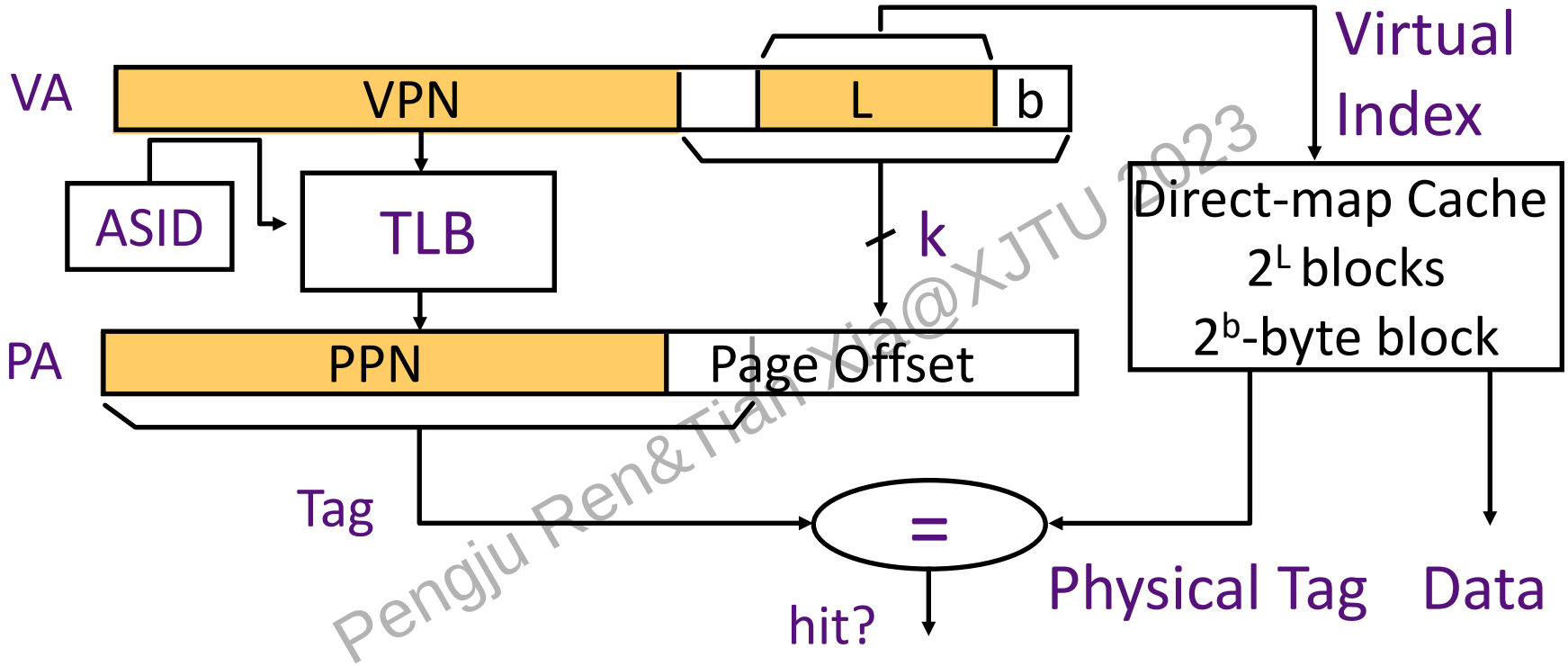
General Solution: *Prevent aliases coexisting in cache*

Software (i.e., OS) solution for direct-mapped cache:

VAs of shared pages must **agree in cache index bits**; this ensures all VAs accessing same PA will conflict in direct-mapped cache (early SPARCs)

Concurrent Access to TLB & Cache
(Virtual Index/Physical Tag, VIPT)

L=Cache Index, b=Cache Block, k=Page Size,



Tag comparison is made after both accesses are completed.

if Cases: $\underline{L + b \le k}$

Index L is available without consulting the TLB.

→ Cache and TLB accesses can begin simultaneously!

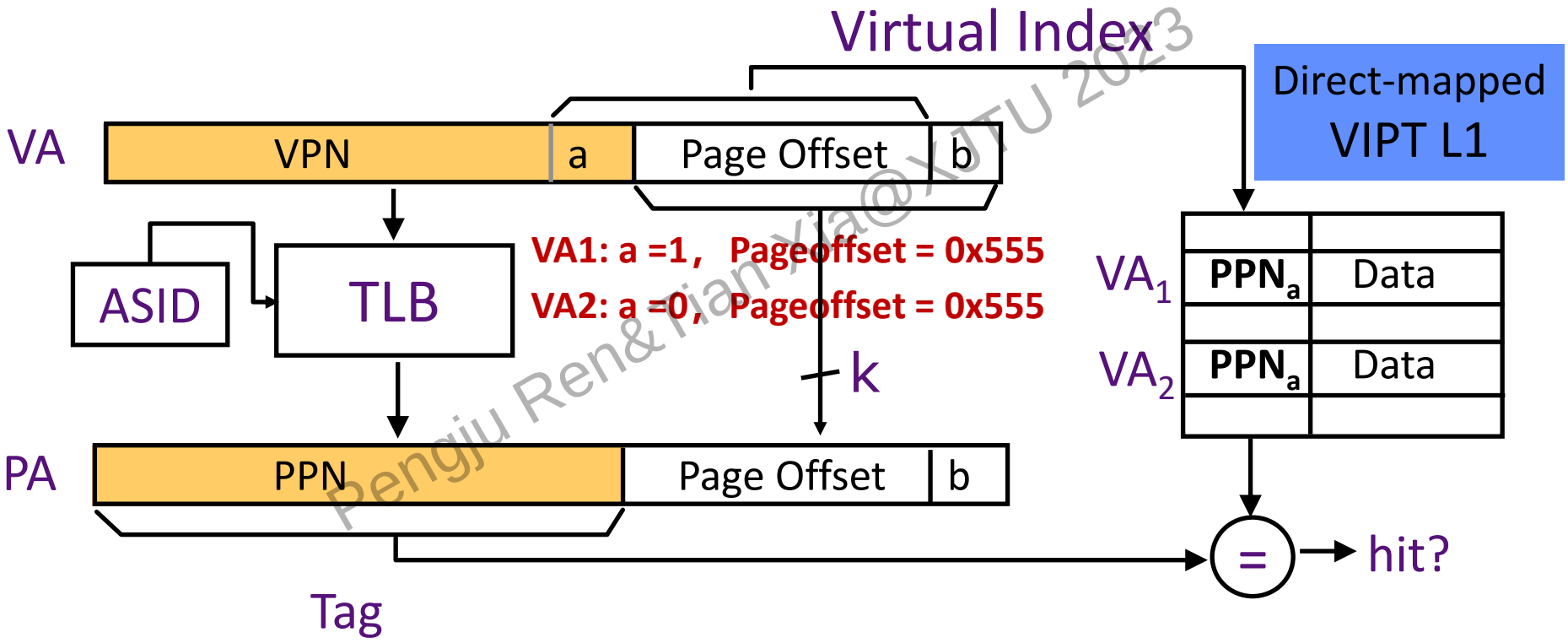
→ Actually, these are **physical indexed Cache**

Concurrent Access to TLB & Large L1

The problem with L1 Cache > Page size

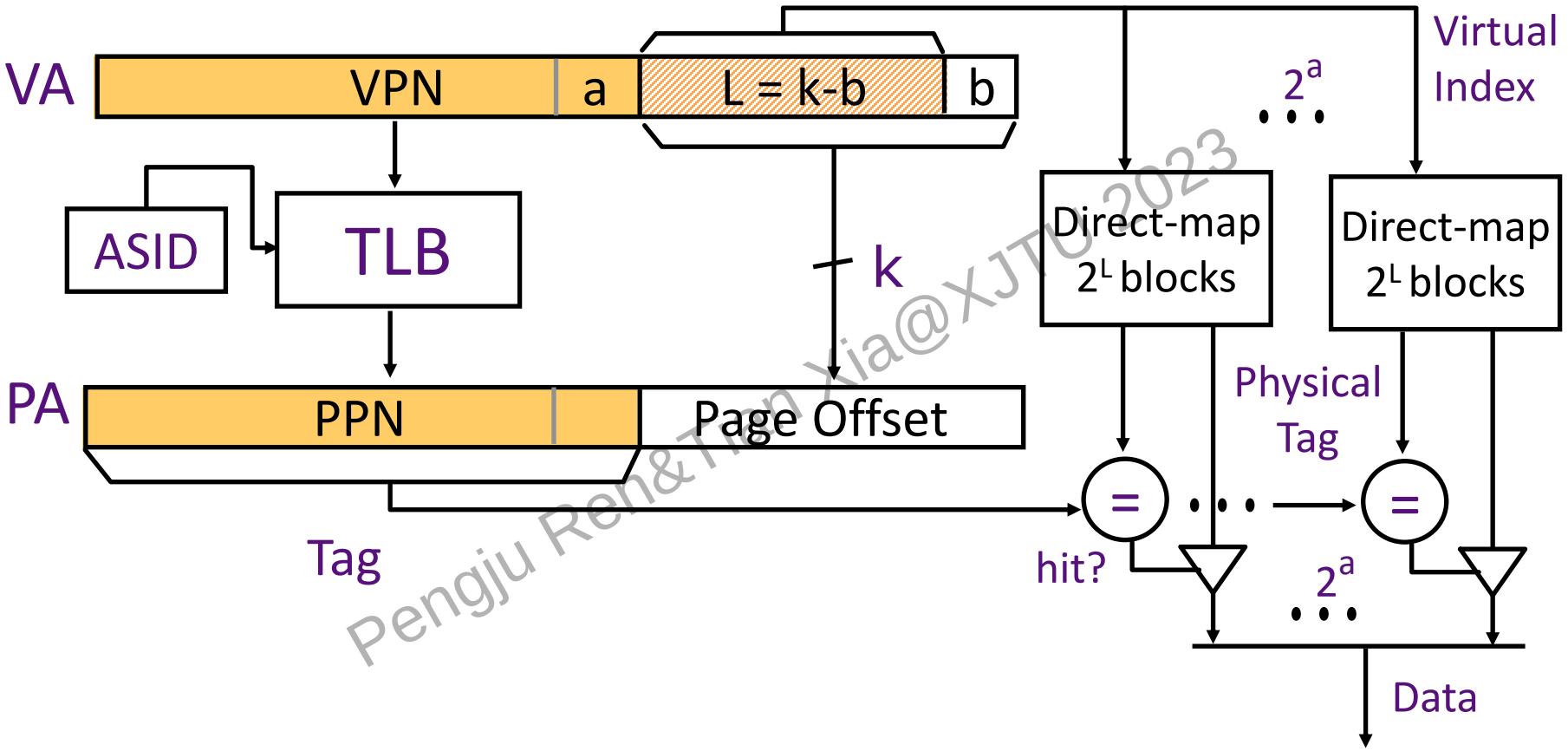
L=Cache Index, b=Cache Block, k=Page Size,

If Cases: L + b > k



VA_1 *and* VA_2 *both map to same* PA

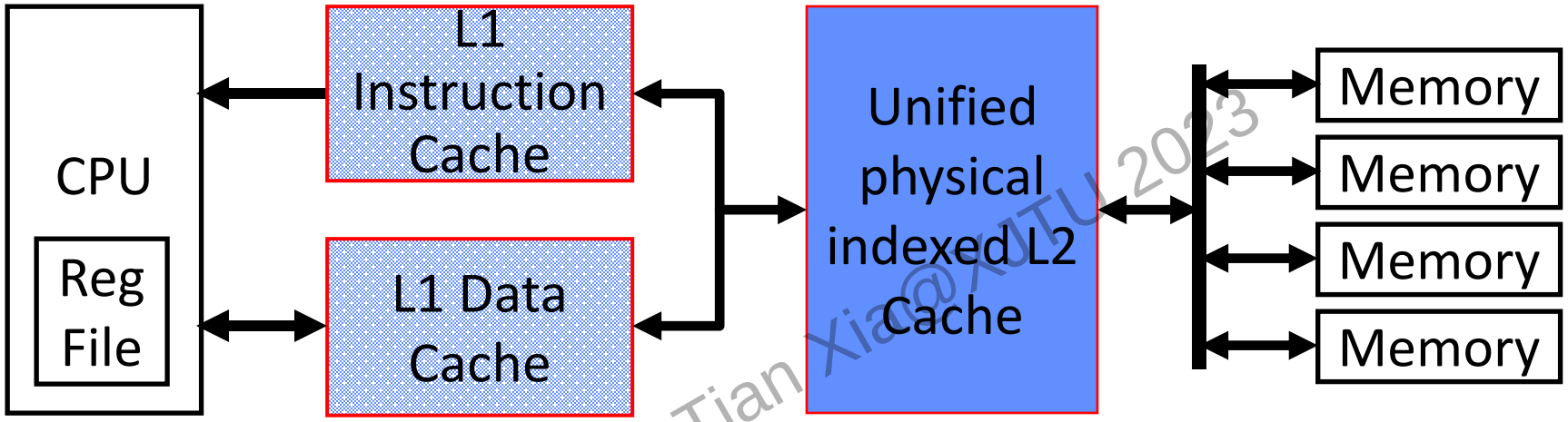
Virtual-Index Physical-Tag Caches:
Associative Organization



After the PPN is known, 2^a physical tags are compared

How does this scheme scale to larger caches?

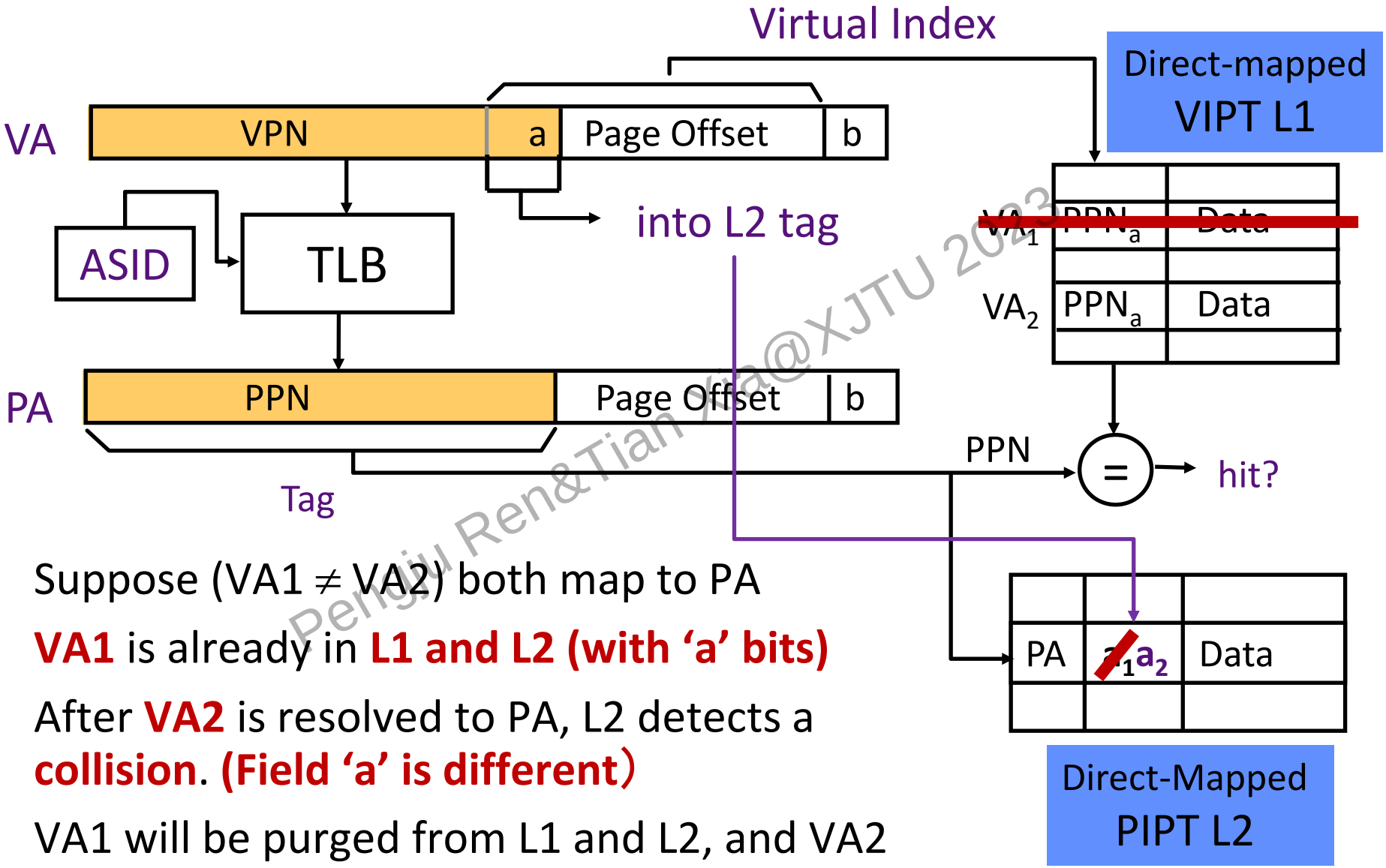
A solution via Second-Level Cache



Physical-index physical-tag (PIPT) , Inclusive L2 Cache:

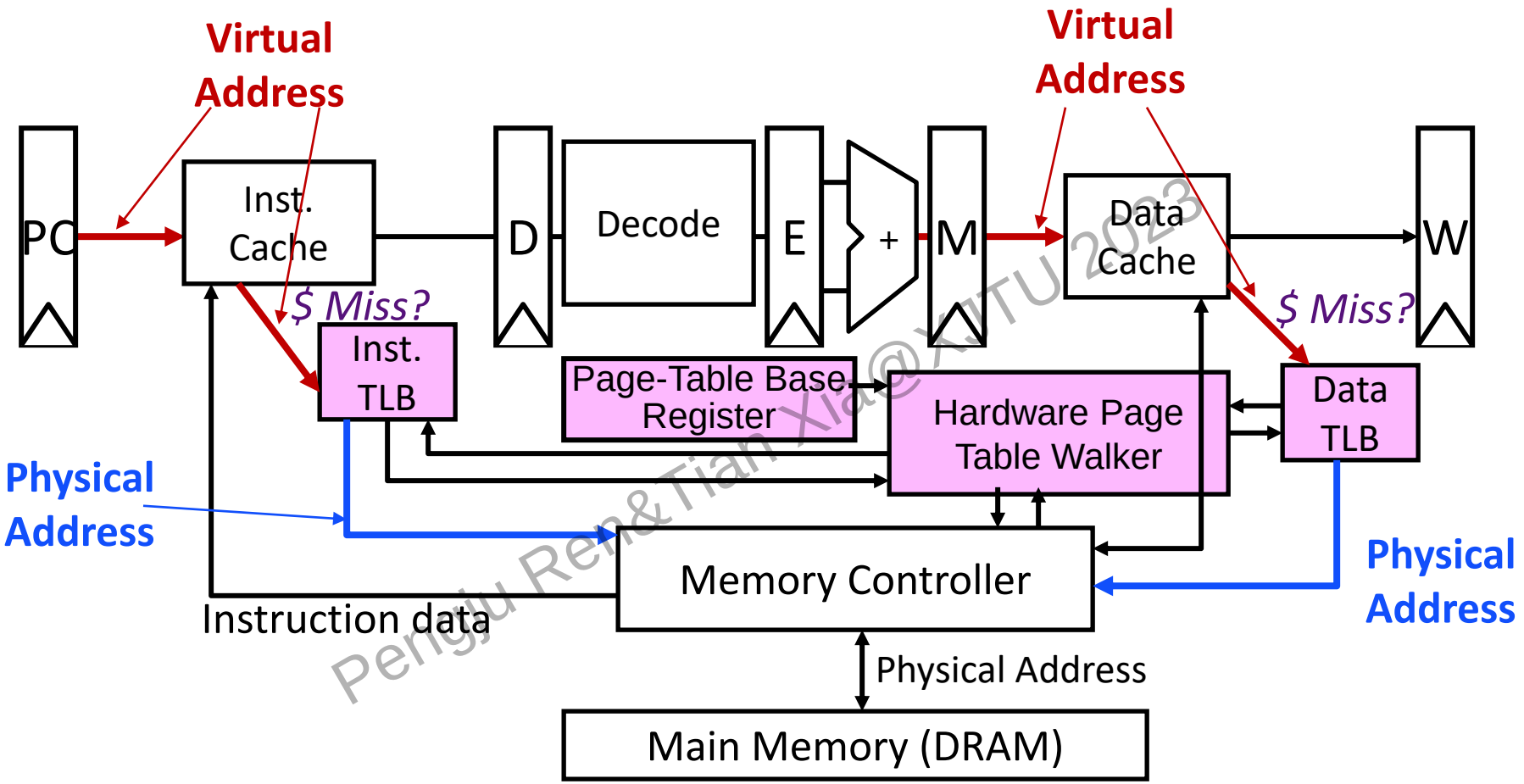
- **PIPT**: Applied for most L2/L3/LLC cache design
- **Unified**: L2 cache backs up both Instruction and Data L1 caches
- **Inclusive:** L2 has copies of all cache lines in both L1 D-Cache and I-Cache

Anti-Aliasing (VIPT) Using L2 [*MIPS R10000,1996*]



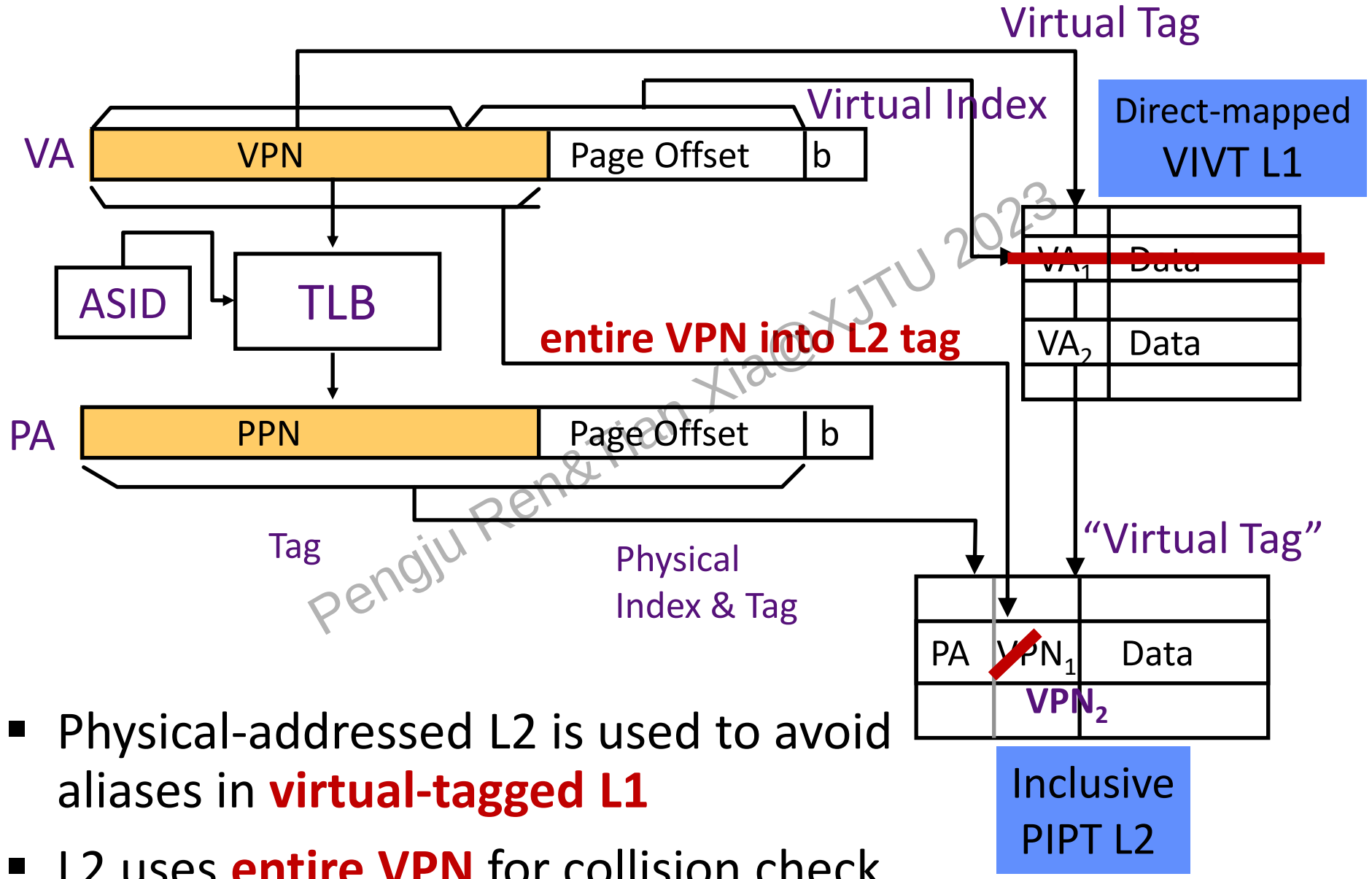
- Suppose (VA1 ≠ VA2) both map to PA
- **VA1** is already in **L1 and L2 (with 'a' bits)**
- After **VA2** is resolved to PA, L2 detects a **collision**. **(Field 'a' is different)**
- VA1 will be purged from L1 and L2, and VA2 will be loaded ⇒ *no aliasing !*

Virtually Addressed Cache
(Virtual Index/Virtual Tag, VIVT)



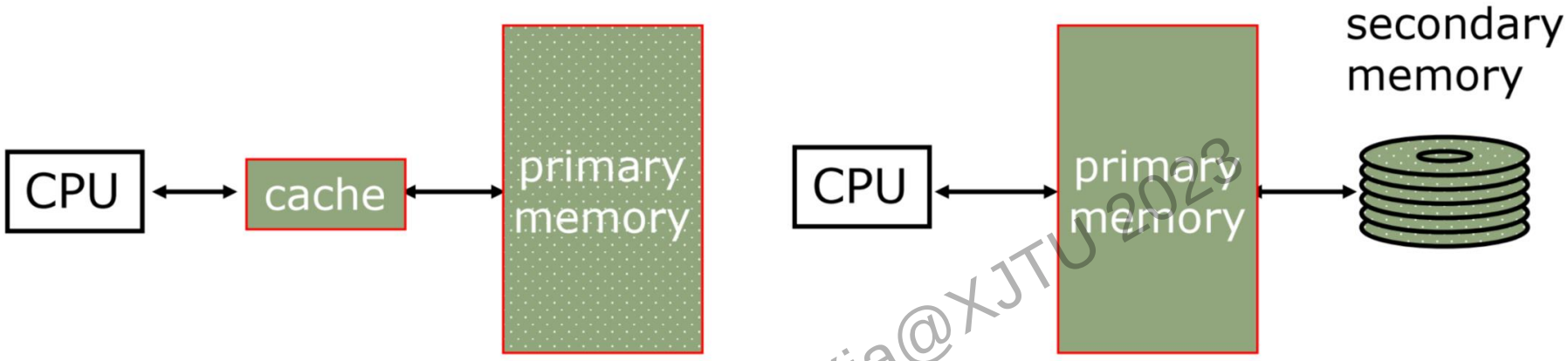
- Directly use **full virtual address** for L1 cache access
- Only check TLB on **L1 cache miss**
- Use **physical address** for L2 cache access

Anti-Aliasing (VIVT) using L2



- Physical-addressed L2 is used to avoid aliases in **virtual-tagged L1**
- L2 uses **entire VPN** for collision check

Summary: Caching v.s Demand Paging



Caching

- cache entry
- cache block (~32 bytes)
- cache miss rate (1% to 20%)
- cache hit (~1 cycle)
- cache miss (~100 cycles)
- a miss is handled
  - in *hardware*

Demand paging

- page frame
- page (~4K bytes)
- page miss rate (<0.001%)
- page hit (~100 cycles)
- page miss (~5M cycles)
- a miss is handled
  - mostly in *software*

Summary: TLB, Cache and Page

TLB、 Page Table、 D-Cache and Physical Page (Total 2^4=16 cases)

- IF PT hit → Physical Page exists in DDR
- IF TLB hit **or** D$ hit → PT hit **and** physical page exists in DDR

Cost ↓

Situation	TLB	Page Table	D-Cache	Physical Page	Operations
D$ hit	**Hit**	**TLB hit, it hit**	**Hit**	**D$ hit, it hit**	No need to check Phys. Mem
D$ miss	**Hit**	**TLB hit, it hit**	**Miss**	**TLB hit, it hit**	Update D$
TLB miss	**Miss**	**D$ hit, it Hit**	**Hit**	**D$ hit, it hit**	Update TLB, then access TLB again
TLB+D$ both miss	**Miss**	**Hit**	**Miss**	**PT hit, it hit**	Update TLB and D$, then access TLB again
Page Fault	**Miss**	**Miss**	**Miss**	**Miss**	Page Fault process

Summary : TLB, Cache and Page

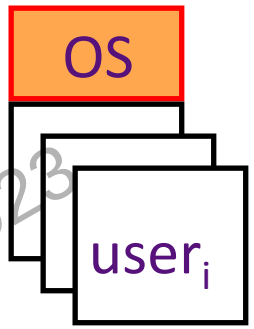
TLB	PAGE TABLE	CACHE	POSSIBLE?	NOTES
Hit	Hit	Miss	Yes	Page Table is never checked if TLB hits
Miss	Hit	Hit	Yes	TLB misses, entry found in Page Table. After retry data is found in Cache
Miss	Hit	Miss	Yes	TLB misses, entry found in Page Table. After retry data misses in Cache
Miss	Miss	Miss	Yes	TLB misses, followed by a Page Fault. After retry data misses in Cache
Hit	Miss	Miss	No	Cannot have translation in TLB if Page Table not in memory
Hit	Miss	Hit	No	Cannot have translation in TLB if Page Table not in memory
Miss	Miss	Hit	No	Data cannot be allowed in Cache if Page Table is not memory

Summary：Modern Virtual Memory Systems

Illusion of a large, private, uniform store

Protection & Privacy

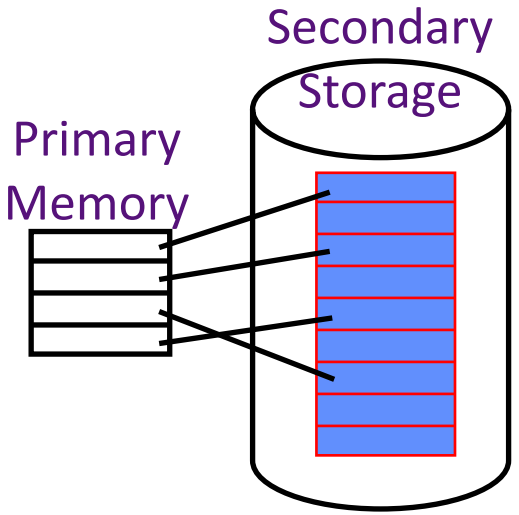
Several users, each with their private address space and one or more shared address spaces



Demand Paging

Provides the ability to run programs larger than the primary memory

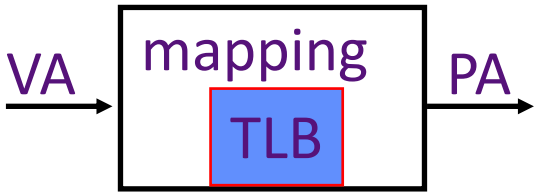
Hides differences in physical memory layouts



Cost

The price is address translation on each memory reference

Use TLB and Virtual Indexed cache



Why a Privileged Architecture?

- Profiles (Simple Embedded w/wo Protection, Unix-like OS, Cloud OS)
- Privileges and Modes
- Privileged Features
  - CSRs
  - Instructions
- Memory Addressing
  - Translation
  - Protection
- Trap Handling
  - Exceptions
  - Interrupts
- Counters
  - Time
  - Performance

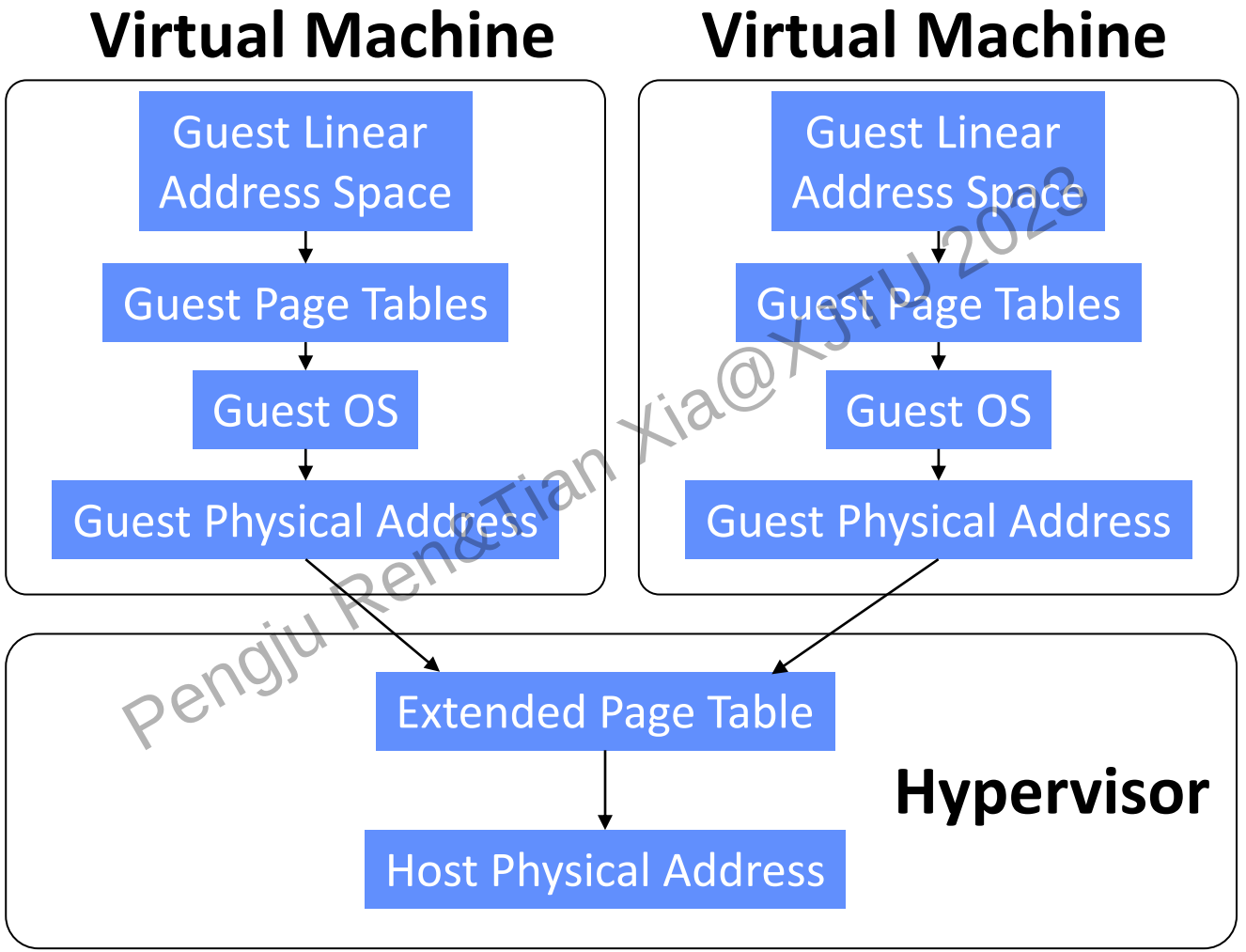
RISC-V Privilege Modes

- Machine mode (M-mode, highest privileges)
  - A.K.A monitor mode, microcode mode, …
- Hypervisor-Extended Supervisor Mode (HS-Mode)
- Supervisor Mode (S-mode)
- User Mode (U-mode, lowest privileges)
- Supported combinations of modes:
  - M (simple embedded systems)
  - M, U (embedded systems with security)
  - M, S, U (systems running Unix-like OS)
  - M, S, HS, U (systems running hypervisors, Cloud OS Capable)

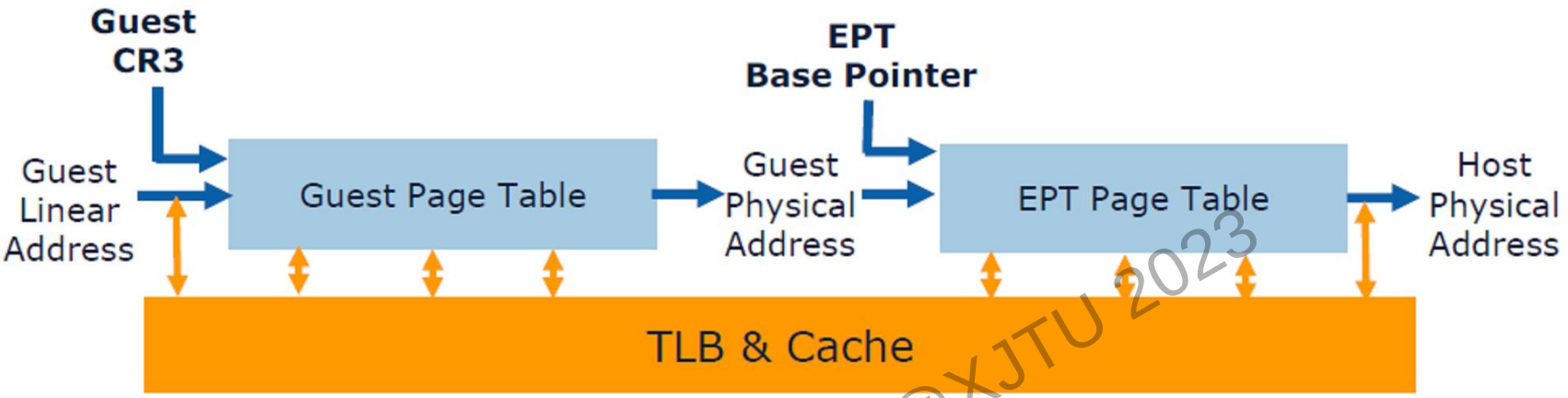
RISC-V System State (Privileged)

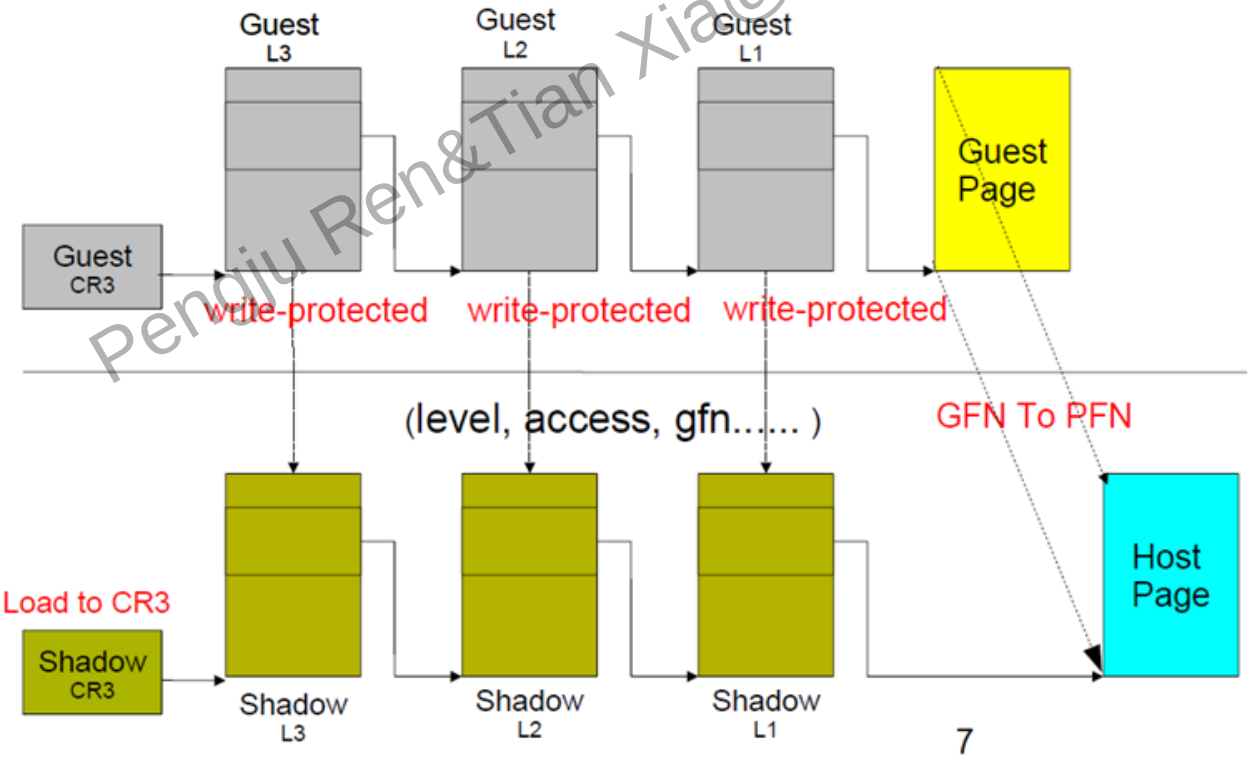
- Processor registers
  - Compute registers
    - General-purpose (x0-x31)
    - Optional floating-point (f0-f31)
    - Optional vector (v0-v31)
    - Optional custom
  - **Control and status registers (CSRs)**
    - **Accessibility controlled by privilege mode or higher**
- System main memory
- System I/O devices
- All system memory and device control registers mapped into flat machine physical address space

Extended Page Tables (EPT) in HS Mode

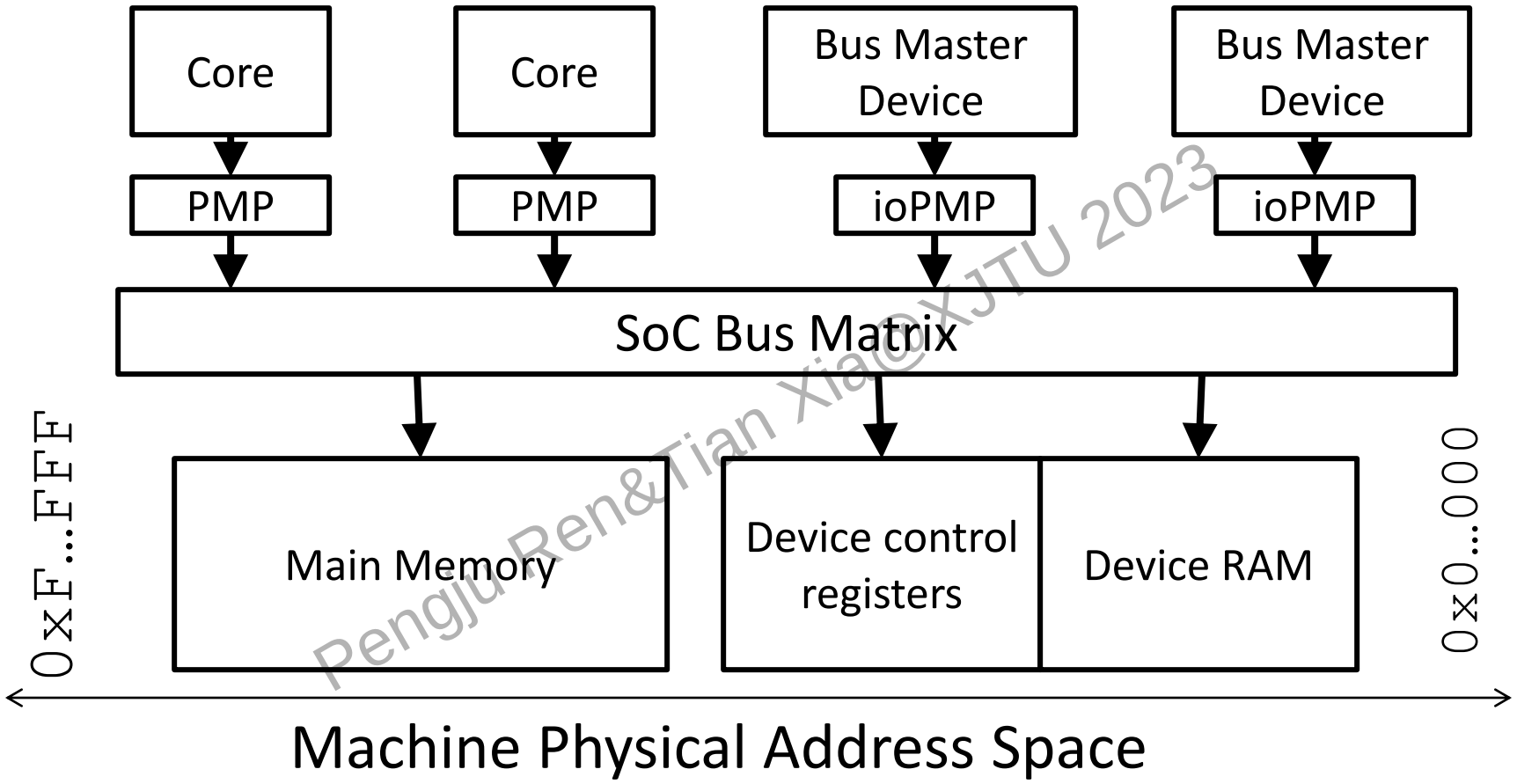


Extended Page Tables (EPT) in HS Mode





Physical Memory Protection (PMP)

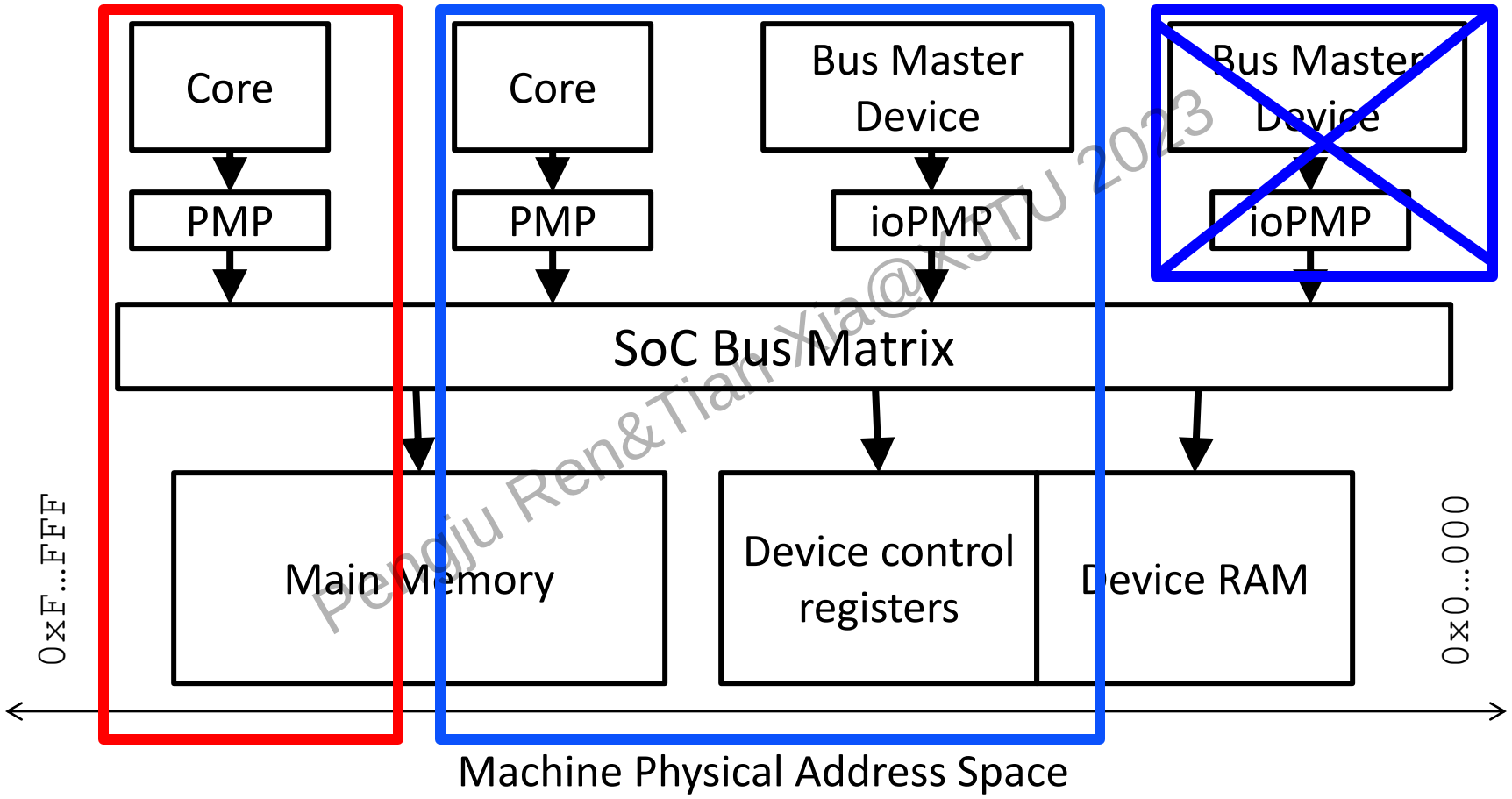


An optional **physical memory protection (PMP)** unit provides per-hart machine-mode control registers to allow physical memory access privileges (read, write, execute) to be specified for each physical memory region

M-Mode controls PMPs

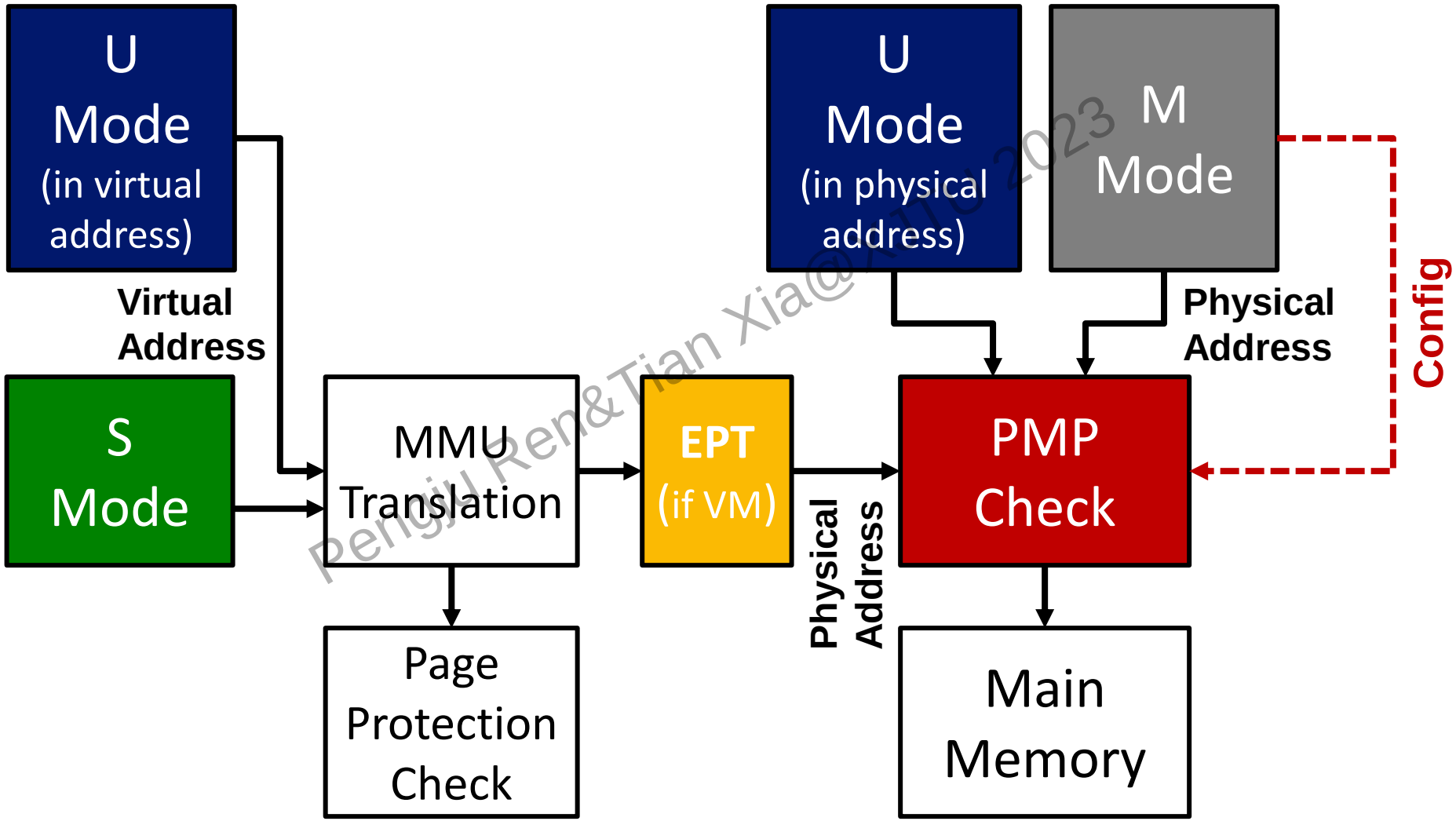
- M-mode has access to entire machine after reset
- Configures PMPs and ioPMPs to contain each active context inside a physical partition
- Can even restrict M-mode access to regions until next reset
- M-mode can dynamically swap PMP settings to run different security contexts on a hart
- RISC-V hardware thread (*hart*)

Multiple Concurrent Security Contexts



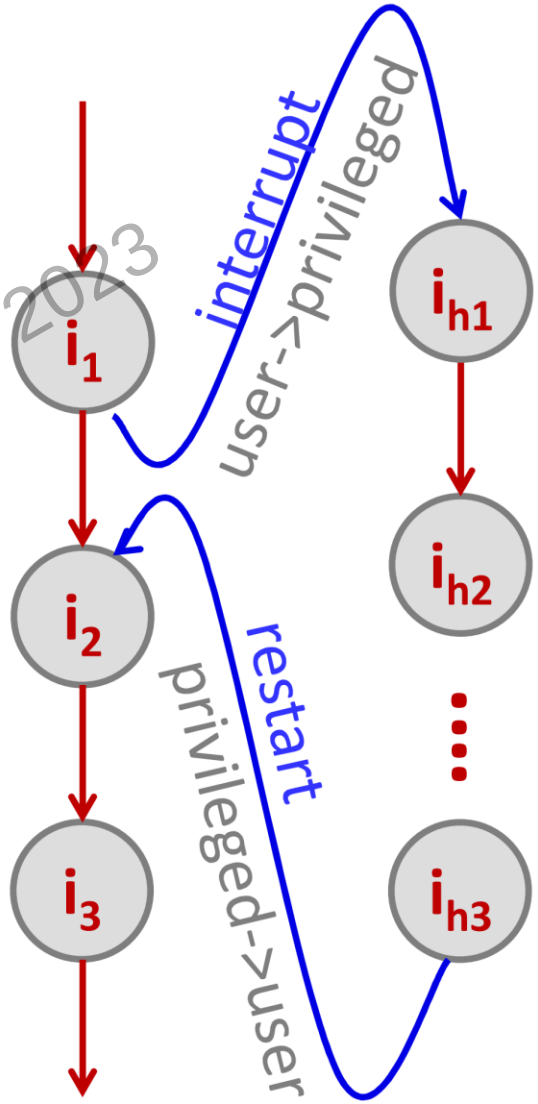
PMP checks are applied to all accesses when the hart is running in S or U modes

Memory Protection for RISC-V Modes



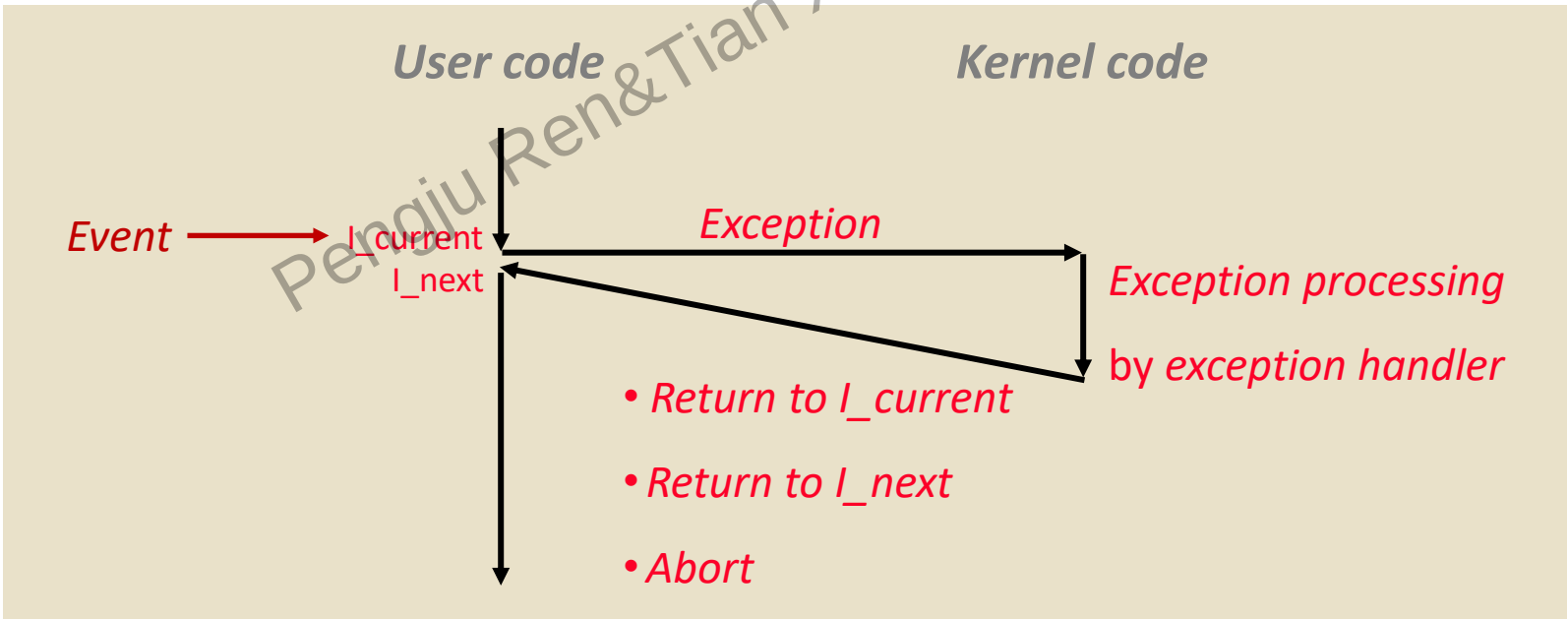
Privilege Level Change

- Combine privilege level change with **interrupt/exception** transfer
  – switch to next higher privilege level on interrupt/exception
  – privilege level restored on return from interrupt/exception
- Interrupt/exception control transfer is **only gateway** to privileged mode
  – lower-level code can never escape into privileged mode
  – lower-level code don't even need to know there is a privileged mode

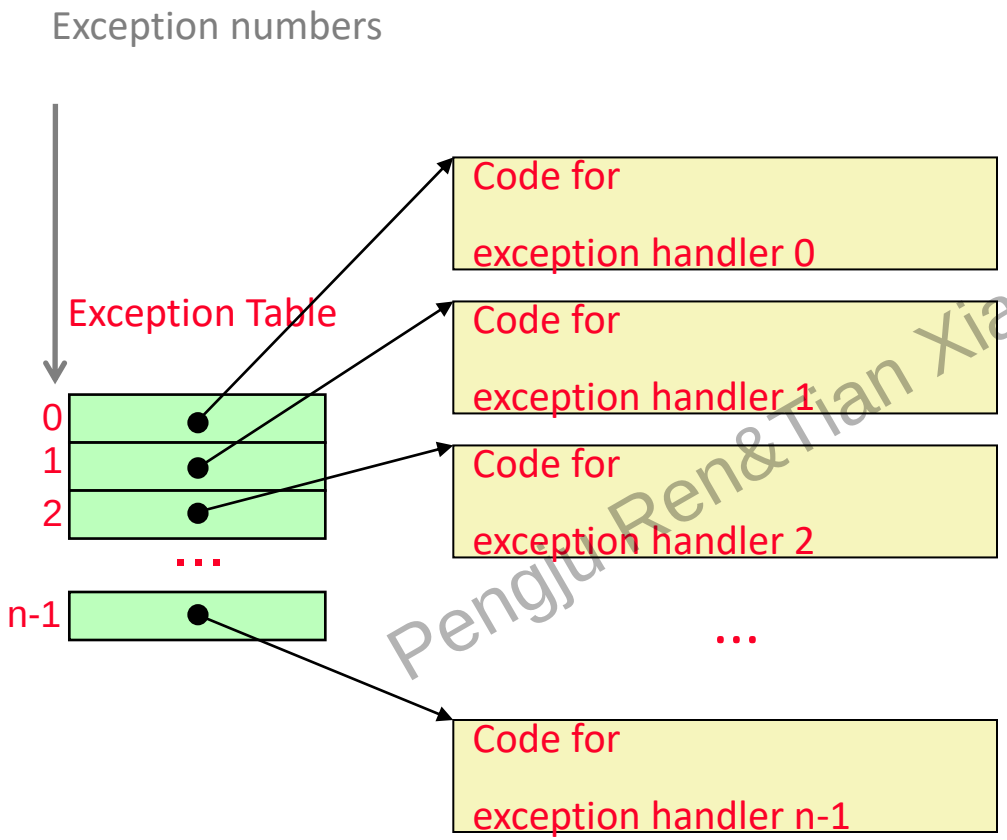


Exceptions

- An ***exception*** is a transfer of control to the OS *kernel* in response to some *event* (i.e., change in processor state)
  - Kernel is the memory-resident part of the OS
  - Examples of events: Divide by 0, arithmetic overflow, page fault, I/O request completes, typing Ctrl-C



Exception Tables



- **Each type of event has a unique exception number k**
- **k = index into exception table (a.k.a. interrupt vector)**
- **Handler k is called each time exception k occurs**

Asynchronous Exceptions (Interrupts)

- Caused by events external to the processor
  - Indicated by setting the processor's *interrupt pin*
  - Handler returns to "**next**" instruction
- Examples:
  - Timer interrupt
    - Every few ms, an external timer chip triggers an interrupt
    - Used by the kernel to take back control from user programs
  - I/O interrupt from external device
    - Hitting Ctrl-C at the keyboard
    - Arrival of a packet from a network
    - Arrival of data from a disk

Synchronous Exceptions

- Caused by events that occur as a result of executing an instruction:
  - ***Traps***
    - Intentional, set program up to "trip the trap" and do something
    - Examples***: software interrupts, system calls***, gdb breakpoints
    - Returns control to "**next**" instruction
  - ***Faults***
    - Unintentional but possibly recoverable
    - Examples: page faults (recoverable), protection faults (unrecoverable), floating point exceptions
    - Either re-executes faulting ("**current**") instruction or aborts
  - ***Aborts***
    - Unintentional and unrecoverable
    - Examples: illegal instruction, parity error, machine check
    - Aborts current program

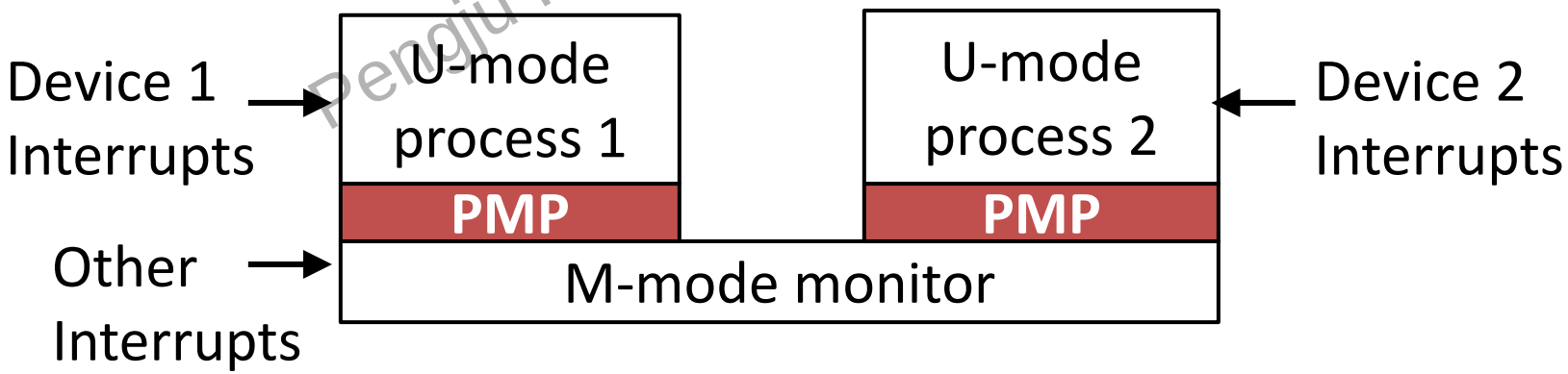
System Calls

- **Each x86-64 system call has a unique ID number**
- **Examples:**

Number	*Name*	*Description*
0	`read`	Read file
1	`write`	Write file
2	`open`	Open file
3	`close`	Close file
4	`stat`	Get info about file
57	`fork`	Create process
59	`execve`	Execute a program
60	`_exit`	Terminate process
62	`kill`	Send signal to process

RISC-V Secure Embedded Systems (M, U modes)

- M-mode runs secure boot and runtime monitor
- Embedded code runs in U-mode
- Physical memory protection (PMP) on U-mode accesses
- Interrupt handling can be delegated to U-mode code
  - User-level interrupt support (N-extension)
- Provides arbitrary number of isolated security contexts

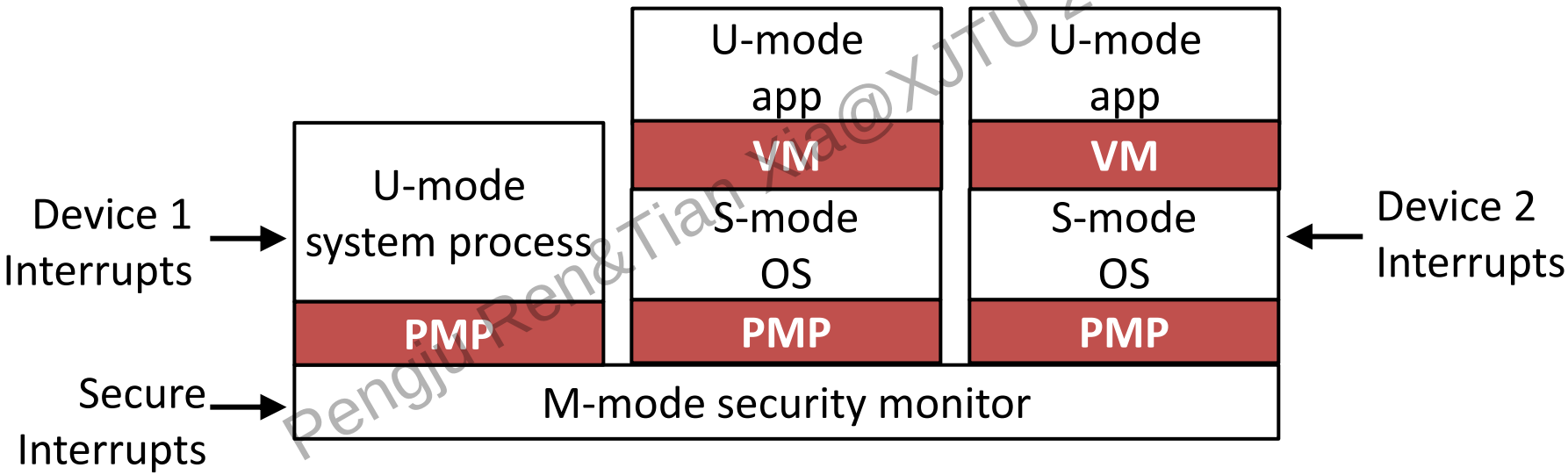


RISC-V Virtual Memory Architectures (M, S, U modes)

- Designed to support current Unix-style operating systems
- Sv32 (RV32)
  - Demand-paged 32-bit virtual-address spaces
  - 2-level page table
  - 4 KiB pages, 4 MiB megapages
- Sv39 (RV64)
  - Demand-paged 39-bit virtual-address spaces
  - 3-level page table
  - 4 KiB pages, 2 MiB megapages, 1 GiB gigapages
- Sv48, Sv57, Sv64 (RV64)
  - Sv39 + 1/2/3 more page-table levels

S-Mode runs on top of M-mode

- M-mode runs secure boot and monitor
- S-mode runs OS
- U-mode runs application on top of OS or M-mode



- PMP checks are also applied to page-table accesses for virtual-address translation, for which the effective privilege mode is S. Optionally, PMP checks may additionally apply to M-mode accesses.
- PMP can *grant* permissions to S and U modes, which by default have none, and can *revoke* permissions from M-mode, which by default has full permissions.

Virtual Memory Use Today - 1

- Servers/desktops/laptops/smartphones have full demand-paged virtual memory
  - Portability between machines with different memory sizes
  - Protection between multiple users or multiple tasks
  - Share small physical memory among active tasks
  - Simplifies implementation of some OS features
- Vector supercomputers have translation and protection but **rarely complete demand-paging**
- (Older Crays: base&bound, Japanese & Cray X1/X2: pages)
  - **Don't waste expensive CPU time** thrashing to disk (make jobs fit in memory)
  - Mostly run in batch mode (run set of jobs that fits in memory)
  - Difficult to implement restartable vector instructions

Virtual Memory Use Today - 2

- Most embedded processors and DSPs provide **physical addressing only**
  - Can't afford area/speed/power **budget** for virtual memory support
  - Often there is no secondary storage to swap to!
  - Programs custom written for particular memory configuration in product
  - Difficult to implement restartable instructions for exposed architectures

Next Lecture： Branch Prediction

Acknowledgements

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